

Mathematical Geometry Processing: Laplacian and Beyond

Yu Wang

Harvard Medical School

MIT Computer Science & Artificial Intelligence Laboratory



Geometry, Diffeomorphism, Inverse PDE & Operator Learning



Joint work with

- Prof. Mirela Ben-Chen, Technion
- Prof. Iosif Polterovich, U. of Montreal
- Dr. Vladimir Kim, Adobe Research
- Prof. Michael Bronstein, Oxford
- Minghao Guo, MIT
- Prof. Justin Solomon, MIT



Review: Operators for Geometric Computing

"Intrinsic and Extrinsic Operators for Shape Analysis" "Variational Quasi-Harmonic Maps for Computing Diffeomorphisms"

Machine Learning on 3D Shapes?

• Voxel grid



• Point cloud



PointNet [Qi et al. 2016]

• Mesh





[Maron et al. 2017]



Geodesic CNN [Masci et al. 2015]

3D ShapeNets [Wu et al. 2015]

The Operator/Matrix for Geometric Computing

• Geometry *representation*: many possibilities!





triangle meshes, polygonal meshes,

triangle soups, implicit representations...

- distance on point clouds [Crane et al. 2013]
 - Q: algorithms behave *consistently* across representations?
 - A: work with continuous operator and discretize it as a matrix.
 - geometric computing: identify the right computational model.

Bridging Geometric Computing and Applied Mathematics



- identify / discover mathematics most relevant in
- designing geometric algorithms and numerical optimization
- that perform best on applications with empirical evaluation metrics
- math/geometry ideas → discretized PDEs/discrete representation → optimization & numerical algorithms → geometric algorithms → applications

Non-Euclidean Signals: Laplacian Spectral Basis

• Euclidean domain \mathbb{T}^d : Fourier bases, 1D: $\sin(nx)$, $\cos(nx)$

•
$$\Delta = \nabla \cdot \nabla = \sum_{i=1}^{d} \frac{\partial^2}{\partial x_i^2} \qquad \Delta(\sin(nx)) = -n^2 \sin(nx)$$

• Spherical domain: spherical harmonics

• $\Delta = ...$



• Any non-Euclidean domain: eigenfunctions of Laplace-Beltrami (Laplacian)



$$\Delta \phi_i = \lambda \phi_i$$

Laplacian eigenfunctions [Levy 2006]

Laplacian Operator on Meshes?

$$\Delta = \nabla \cdot \nabla = \frac{\partial^2}{\partial x_1^2} + \frac{\partial^2}{\partial x_2^2}$$

Image Processing.



- Laplacian on images
 - finite difference as approximation

•
$$\frac{df(x)}{dx} \approx \frac{f(x+1) - f(x-1)}{2}$$

•
$$\Delta f(x,y) \approx \frac{f(x+1,y) + f(x-1,y) + f(x,y+1) + f(x,y-1) - 4 \times f(x,y)}{f(x+1,y) + f(x-1,y) + f(x,y+1) + f(x,y-1) - 4 \times f(x,y)}$$

4

• via the stencil:

$$\begin{bmatrix} 0 & 1 & 0 \\ 1 & -4 & 1 \\ 0 & 1 & 0 \end{bmatrix}$$

Mesh Processing.



- Laplacian on meshes?
 - similar idea, $\Delta u \approx \sum_{j \in \mathcal{N}(i)} \mathcal{L}_{ij}(\mathbf{u}_j \mathbf{u}_i)$ but
 - *irregular connectivity*
 - entries L_{ij} depends on the shape of triangles



Triangle Mesh as Graph

- Graph: nodes connected by edges (${\cal E}$)
 - Dirichlet energy for $u \in \mathbb{R}^{n \times 1}$ (#nodes by 1)

$$E(\mathbf{u}) = \frac{1}{2} \sum_{\{i,j\} \in \mathcal{E}} w_{ij} (\mathbf{u}_i - \mathbf{u}_j)^2 = \frac{1}{2} \mathbf{u}^{\mathrm{T}} \mathbf{L} \mathbf{u}$$

- adjacency matrix $J \in \mathbb{R}^{e \times n}$ (#edges by #nodes)
- graph Laplacian $L \in \mathbb{R}^{n \times n}$ (#nodes by #nodes)

$$L_{ij} = \begin{cases} w_{ij} & \text{if } \{i, j\} \text{ is an edge} \\ -\sum_{\substack{j \neq i \\ 0}} L_{ij} & \text{if } i = j \\ 0 & \text{otherwise} \end{cases} \quad J_{ei} = \begin{cases} 1 & \text{if } e \text{ is an edge starts at } i \\ -1 & \text{if } e \text{ is an edge ends at } i \\ 0 & \text{otherwise} \end{cases}$$
$$L = J^{T} \text{diag(w) } J$$

- Mesh: still a graph, except the edges comes from triangles
 - graph has to be embeddable in 3D/2D
 - edges are associated with some lengths (with triangle inequality)

9

Triangle Mesh as Graph

- Mesh: vertices connected by triangles (\mathcal{T})
 - Dirichlet energy for $u \in \mathbb{R}^{n \times 1}$ (#nodes by 1)

$$E(\mathbf{u}) = \frac{1}{2} \sum_{t = \{i, j, k\} \in \mathcal{T}} ||\mathbf{G}_{ti}\mathbf{u}_i + \mathbf{G}_{tj}\mathbf{u}_j + \mathbf{G}_{tk}\mathbf{u}_k||^2 = \frac{1}{2} \mathbf{u}^{\mathrm{T}} \mathbf{L} \mathbf{u}$$

- grad matrix $\mathbf{G} \in \mathbb{R}^{2f \times n}$ (2#faces by #vertices)
 - assume u is a piece-wise linear function in the triangle $t = \{i, j, k\}$
 - $\nabla u = (G_{ti}u_i + G_{tj}u_j + G_{tk}u_k) \in \mathbb{R}^2$ is the grad of u in $t = \{i, j, k\}$
- "mesh" Laplacian L: still a graph Laplacian
 - with weights depends on local geometry, i.e., G

$$\mathbf{L}_{ij} = \begin{cases} w_{ij} & \text{if } \{i, j\} \text{ is an edge} \\ -\sum_{j \neq i} \mathbf{L}_{ij} & \text{if } i = j \\ 0 & \text{otherwise} \end{cases}$$

$$\mathbf{L} = \mathbf{G}^{\mathrm{T}}\mathbf{G}$$



- Graph Laplacian $\mathbf{L} = \mathbf{J}^{\mathrm{T}}\mathbf{J}$
 - J: edge-node adjacency matrix
 - $J \in \mathbb{R}^{e \times n}$ (#edges by #nodes)
 - Ju: the difference per edge for $u \in \mathbb{R}^{n \times 1}$

- Mesh Laplacian $\mathbf{L} = \mathbf{G}^{\mathrm{T}}\mathbf{G}$
 - G: face-vertex gradient matrix
 - $G \in \mathbb{R}^{2f \times n}$ (2#faces by #vertices)
 - Gu: the gradient per triangle for $\mathbf{u} \in \mathbb{R}^{n \times 1}$

Harmonic Functions: Mean Value Property

- Optimization: $\min \frac{1}{2}u^{T}Lu \text{ s.t. } R^{T}u = b$
 - $R^T u = b$: the constraint that u is known at some nodes
 - R is the binary selection matrix choosing known rows in \boldsymbol{u}
 - S is the binary selection matrix choosing unknown rows in u
 - $S^T R = 0$
- Lagrangian multiplier method:
 - min $\frac{1}{2}u^{T}Lu + \lambda^{T}(R^{T}u b)$
 - $Lu + R\lambda = 0 \rightarrow S^{T}Lu = 0$
- This implies the unknown \mathbf{u}_i is

harmonic: u_i equals to the weighted average of its neighbors' values

$$\mathbf{u}_{i} = \frac{1}{\sum_{j \in \mathcal{N}(i)} \mathbf{L}_{ij}} \sum_{j \in \mathcal{N}(i)} \mathbf{L}_{ij} \mathbf{u}_{j}$$



Discrete v.s. Continuous: Harmonic Interpolation

• A quadratic optimization with linear constraints minimize the derivation of u_i or u(x) locally

• $\min_{\mathbf{u}\in\mathbb{R}^n} \frac{1}{2} \mathbf{u}^{\mathrm{T}} \mathbf{L} \mathbf{u} \text{ s.t. } \mathbf{R}^{\mathrm{T}} \mathbf{u} = \mathbf{b}$



- $S^T G^T G u = 0$
 - S selects unknown rows in u
 - solve a linear system

• $\min_{u(\mathbf{x})} \frac{1}{2} \int_{\Omega} |\nabla u(\mathbf{x})|^2 d\mathbf{x} \text{ s.t. } u|_{\partial\Omega} = b(\mathbf{x})$





- for $\mathbf{x} \in \Omega \setminus \partial \Omega$: where $u(\mathbf{x})$ is unknown
- solve a linear PDE

• Mean Value Property: u_i or u(x) equals to the weighted average of its neighbors' values

Laplace Equation: PDE 101

Definition (Laplace Equation: Mixed boundary condition)

Consider a partition of the boundary: $\partial \Omega = \partial \Omega_N \cup \partial \Omega_D$ such that $\partial \Omega_N \cap \partial \Omega_D = \emptyset$.

$$\begin{split} \Delta f(\mathbf{x}) &= 0, & \forall \mathbf{x} \in \Omega \setminus \partial \Omega \quad (\text{interior}) \\ f(\mathbf{x}) &= f_0(\mathbf{x}), & \forall \mathbf{x} \in \partial \Omega_D \quad (\text{Dirichlet condition}) \\ \mathbf{n} \cdot \nabla f(\mathbf{x}) &= g_0(\mathbf{x}), & \forall \mathbf{x} \in \partial \Omega_N \quad (\text{Neumann condition}) \end{split}$$







• Fact: forward PDE allows only one boundary condition

A Closely Related Concept: Dirichlet-to-Neumann (DtN) Operator \mathcal{S}

Consider a shape Ω bounded by the surface $\Gamma = \partial \Omega$.

$$\begin{cases} \Delta u(\mathbf{x}) = 0 & \mathbf{x} \in \Omega \\ u(\mathbf{x}) = g(\mathbf{x}) & \mathbf{x} \in \partial \Omega \end{cases}$$

where $g(\Gamma)$ is Dirichlet data

Neumann data
$$g_n = \frac{\partial}{\partial n} u(\Gamma)$$

Dirichlet-to-Neumann (DtN) operator: $\mathcal{S} := g \mapsto g_n$ a.k.a the Steklov-Poincaré operator. (temperature-to-flux, voltage-to-current)



Discrete Laplacian: A Sparse Matrix

continuous
$$\Delta =
abla \cdot
abla$$

discrete $\mathbf{L} = \mathbf{G}^{\mathsf{T}}\mathbf{G} \in \mathbb{R}^{n \times n}$

 $G \in \mathbb{R}^{2f \times n}$: gradient operator (weighted) n: #vertices

f: #faces





$$L_{ij} = \begin{cases} 1/2 \left(\cot \alpha_{ij} + \cot \beta_{ij} \right) & \text{if } \{i, j\} \text{ is an edge} \\ -\sum_{\substack{j \neq i \\ 0}} L_{ij} & \text{if } i = j \\ 0 & \text{otherwise} \end{cases}$$

L is a graph Laplacian with geometry-determined edge weights. Same formula for curved and flat surfaces Finite Element Method (FEM) [Steinbach 2007] Discrete Exterior Calculus (DEC) [Desbrun et al., 2005]



• Distance



- Distance
- Parameterization



- Distance
- Parameterization
- Shape descriptor
- Correspondence





Query: Cat

- Distance
- Parameterization
- Shape descriptor
- Correspondence
- Shape classification

```
1 def Geometry_Processing(mesh):
2 L = Laplacian_Operator(mesh)
3 del mesh # delete the mesh
4 results = Linear_Solve(L)
5 ...
6 return results
```

- Distance [Crane et al. 2013]
- Parameterization [Mullen et al. 2008]
- Shape description [Sun et al. 2009]
- Correspondence [Ovsjanikov et al. 2012]
- Shape classification [Bronstein et al. 2011]
- Shape exploration [Rustamov et al. 2013]
- Deformation [Boscaini et al. 2015]
- Shape optimization...
- Mesh generation...

$$L \rightarrow L^{(A)}, \tilde{L}, \text{ or } S$$

This talk covers:

- Search for a matrix $L^{(A)} = G^{T}AG$: same sparsity pattern to $L = G^{T}G$ (major focus)
- Learn from data FEM kernels to assemble entries \tilde{L}_{ij}
- Design explicitly a different matrix S: more informative/robust

Why?

- New operators \rightarrow (much) more expressive computational models
- Systematically improve potentially every task in geometric computing

Optimization in the Space of Laplacians

"Harmonicity": The Key Notion of "Smoothness"





- Key: what is a smooth function on the mesh/graph?
- Harmonic function: a function **u** whose value at each node/vertex *i* equals to the average over $\mathcal{N}(i)$, the neighbors of *i*

$$\mathbf{u}_{i} = \frac{1}{\sum_{j \in \mathcal{N}(i)} 1} \sum_{j \in \mathcal{N}(i)} \mathbf{u}_{j}$$

$$(\Delta \mathbf{u})_i := \frac{1}{\sum_{j \in \mathcal{N}(i)} 1} \sum_{j \in \mathcal{N}(i)} (\mathbf{u}_i - \mathbf{u}_j)$$

"Harmonicity": The Key Notion of "Smoothness"





- Key: what is a smooth function on the mesh/graph?
- Quasi-harmonic function: a function **u** whose value at each node/vertex i equals to the weighted average over $\mathcal{N}(i)$, the neighbors of i

$$\mathbf{u}_{i} = \frac{1}{\sum_{j \in \mathcal{N}(i)} \mathbf{L}_{ij}} \sum_{j \in \mathcal{N}(i)} \mathbf{L}_{ij} \mathbf{u}_{j}$$

$$\left(\Delta^{(w)}\mathbf{u}\right)_{i} = \frac{1}{\sum_{j \in \mathcal{N}(i)} \mathbf{L}_{ij}} \sum_{j \in \mathcal{N}(i)} \mathbf{L}_{ij} (\mathbf{u}_{i} - \mathbf{u}_{j})$$

Quasi-Harmonic Function and Generalized Laplacians

• Smooth
$$\Delta =
abla \cdot
abla$$

- Harmonic: $\Delta u = 0$
 - u_i equals to the *average* over *i*'s neighbors

 $0 = \sum_{j \in N(i)} (u_j - u_i)$

• Smooth
$$\Delta^{(\mathbf{A})} = \nabla \cdot [\mathbf{A}(\mathbf{x})\nabla]$$

- Quasi-harmonic: $\Delta^{(A)}u = 0$
- *u_i* equals to the *weighted-average* over *i*'s neighbors

$$0 = \sum_{j \in N(i)} \mathsf{L}_{ij}(u_j - u_i)$$



Inverse Problems of PDEs for Computing Diffeomorphisms

"Variational Quasi-Harmonic Maps for Computing Diffeomorphisms." Yu Wang, Minghao Guo, and Justin Solomon. *ACM Transactions on Graphics (TOG) 42(4). ACM SIGGRAPH 2023 Journal Track*.

Diffeomorphisms

- Diffeomorphism ϕ : a *smooth* map with *smooth* inverse (ϕ^{-1} must exist)
 - diffeomorphisms: all physically possible deformation (no negative volume)
- Homeomorphism ϕ : smooth \rightarrow continuous
 - injective: $\phi(x) \neq \phi(y)$ for $x \neq y$
 - inversion-free: det $D\phi(x) > 0$, $\forall x$, positive Jacobian $D\phi(x) \in \mathbb{R}^{2 \times 2}$



Diffeomorphism = Smooth Injective / Inversion-free Mapping

• The map $\phi = (u, v)$ can be a:

deformation, shape representation, correspondence, parameterization, ...



- Foundational, wherever using computers to represent shapes in physics, engineering, shape optimization, computer vision, mesh generation...
- Homeomorphism = Inversion-freeness (Under conditions) e.g. [Lipman 2014] det $[\nabla u \nabla v] > 0$

Our Solution: (Quasi) Harmonic Maps

• Previous works, at high-level:

min E(u, v) s.t. det $[\nabla u \nabla v] > 0$

- largely relying on constrained numerical optimization
- solved by customized barrier / interior point methods
- while minimizing some energy *E*, e.g., the Winslow functional from physics
- Starting point: quasi-harmonic map $\nabla \cdot [\mathbf{A}(\mathbf{x}) [\nabla(u, v)]] = (0, 0)$
- Our method

min R(u, v, A) s.t. $\nabla \cdot [A(x) [\nabla(u, v)]] = (0, 0)$

- First review relevant ideas from
 - geometric graph theory
 - complex analysis / 2D PDEs

Problem: Flatten a Surface subject to Positional Constraints

Problem Setup

Suppose Ω is a two-dimensional Riemannian manifold with disk topology, and consider a planar domain $\Gamma \subset \mathbb{R}^2$ whose boundary $\partial \Gamma$ is a simple closed curve. Assume $\phi = (u, v) : \Omega \to \Gamma$ diffeomorphically maps $\partial \Omega$ onto $\partial \Gamma$. Denote the (given) boundary map as $[b_u, b_v](\cdot) : \partial \Omega \to \mathbb{R}^2$, and denote the outward normal to $\partial \Gamma$ as $\hat{n}(\cdot) : \partial \Omega \to S^1$.



Question?

How to find a map ϕ that is diffeomorphic (and minimizes certain functional)?

Review: Tutte Embedding = Discrete Quasi-Harmonic Maps

- Fixed boundary, interior nodes placed at neighbors' weighted average
 - interior positions found by solving a linear system
- Convex boundary → edges do not intersect (injective!)
 - [Tutte 1962]: created the field of geometric graph theory

HOW TO DRAW A GRAPH

By W. T. TUTTE

[Received 22 May 1962]

1. Introduction

WE use the definitions of (11). However, in deference to some recent attempts to unify the terminology of graph theory we replace the term 'circuit' by 'polygon', and 'degree' by 'valency'.

A graph G is 3-connected (nodally 3-connected) if it is simple and non-separable and satisfies the following condition; if G is the union of two proper subgraphs H and K such that $H \cap K$ consists solely of two vertices u and v, then one of H and K is a link-graph (arc-graph) with ends u and v.

It should be noted that the union of two proper subgraphs H and K of G can be the whole of G only if each of H and K includes at least one edge or vertex not belonging to the other. In this paper we are concerned mainly with nodally 3-connected graphs, but a specialization to 3-connected graphs is made in § 12.

In § 3 we discuss conditions for a nodally 3-connected graph to be planar, and in § 5 we discuss conditions for the existence of Kuratowski subgraphs of a given graph. In §§ 6-9 we show how to obtain a convex representation of a nodally 3-connected graph, without Kuratowski subgraphs, by solving a set of linear equations. Some extensions of these results to general graphs, with a proof of Kuratowski's theorem, are given in §§ 10-11. In § 12 we discuss the representation in the plane of a pair of dual graphs, and in § 13 we draw attention to some unsolved problems.



William T. Tutte







[Images from Kyri Pavlou]

Review: Continuous Quasi-Harmonic Maps

- Quasi / A(x)-harmonic: point x placed at neighbors' weighted average
 - $\nabla \cdot [\mathbf{A}(\mathbf{x})\nabla u(\mathbf{x})] = 0, \quad \forall \mathbf{x} \in \Omega \backslash \partial \Omega$
 - $\nabla \cdot [A(\mathbf{x})\nabla \nu(\mathbf{x})] = 0, \quad \forall \mathbf{x} \in \Omega \backslash \partial \Omega$
- Dirichlet boundary condition: fix the boundary

 $u(\mathbf{x}) = b_u(\mathbf{x}), \quad \forall \mathbf{x} \in \partial \Omega$ $v(\mathbf{x}) = b_v(\mathbf{x}), \quad \forall \mathbf{x} \in \partial \Omega$

- Map (u, v) diffeomorphic/injective for convex boundary
 - A(x) = I: by *RKC theorem* in complex analysis
 - $A(x) \neq I$: generalization by [Alessandrini and Nesi 2001]



Tibor Radó

Theorem [Radó-Kneser-Choquet (RKC)]

Harmonic maps onto convex regions are diffeomorphic.

Quasi-Harmonic Maps, aka, Tutte Embedding

- Fix boundary, interior nodes placed at neighbors' weighted average [Tutte 1962]
- Fail for non-convex boundary: flipped triangles (in red color)---our condition fixes it!
- There is a hope with extra conditions [Gortler, Gotsman, Thurston 2006]



Image from [Du et al. 2021]



Main Theory: Diffeomorphism = {Quasi-harmonic} + { Dirichlet & Neumann BCs}

Theorem (Main result: continuous version)

 $\phi = (u, v)$ is a diffeomorphism if and only if there exist (1) a positive definite tensor field $A(\cdot)$ satisfying $\frac{1}{K}I \leq A(x) \leq KI$, (+ smooth conditions) (2) a positive function $s : \partial \Omega \to \mathbb{R}$ that $s(x) \geq S$, for some K, S > 0, such that ϕ is A(x)-harmonic with a special Cauchy boundary condition, i.e.:

> $\nabla \cdot [A(x)\nabla u(x)] = 0 \qquad \forall x \in \Omega \setminus \partial \Omega$ $\nabla \cdot [A(x)\nabla v(x)] = 0 \qquad \forall x \in \Omega \setminus \partial \Omega$ $u(x) = b_u(x) \qquad \forall x \in \partial \Omega$ $v(x) = b_v(x) \qquad \forall x \in \partial \Omega$ $n(x)^{\mathsf{T}} \Big[A(x)\nabla u(x) \quad A(x)\nabla v(x) \Big] = s(x)\hat{n}(x)^{\mathsf{T}} \qquad \forall x \in \partial \Omega$

Main Theory: Diffeomorphism = {Quasi-harmonic} + { Dirichlet & Neumann BCs}

• Main Theorem: map (u, v) is diffeomorphic *if-and-only-if* such an A(x) exists:

• (1) Quasi-harmonic:

Same

Ours

$$\nabla \cdot [A(\mathbf{x})\nabla u(\mathbf{x})] = 0, \quad \forall \mathbf{x} \in \Omega \setminus \partial \Omega$$
$$\nabla \cdot [A(\mathbf{x})\nabla v(\mathbf{x})] = 0, \quad \forall \mathbf{x} \in \Omega \setminus \partial \Omega$$

Tutte • (2) Dirichlet boundary condition: specify the boundary positions (trivial).

$$u(\mathbf{x}) = b_u(\mathbf{x}), \quad \forall \mathbf{x} \in \partial \Omega$$
$$v(\mathbf{x}) = b_v(\mathbf{x}), \quad \forall \mathbf{x} \in \partial \Omega$$

• (3) Neumann boundary condition: specify the A(x)-weighted normal derivative. $n(x)^{\mathsf{T}}[A(x)\nabla u(x) \quad A(x)\nabla v(x)] = \hat{n}(x)^{\mathsf{T}} \quad \forall x \in \partial \Omega$

- **n**(x): normal on source domain
- $\widehat{\boldsymbol{n}}(\boldsymbol{x})$: normal on target domain
A Condition *Feasible* Computationally

Neumann boundary condition: specify the A(x)-weighted normal derivative $n(x)^{T}[A(x)\nabla u(x) \quad A(x)\nabla v(x)] = \hat{n}(x)^{T} \quad \forall x \in \partial \Omega$

- **n**(x): normal on source domain
- $\widehat{\mathbf{n}}(\mathbf{x})$: normal on target domain

Ours: correct & *feasible* computationally

the *nature* boundary condition: the best thing you can hope for! made possible *discrete* injectivity

$$\mathbf{n}(\mathbf{x})^{\mathsf{T}} [\nabla u(\mathbf{x}) \quad \nabla v(\mathbf{x})] = \hat{\mathbf{n}}(\mathbf{x})^{\mathsf{T}} \quad \forall \mathbf{x} \in \partial \Omega$$

A possible variant:

theoretically correct

computationally *inconsistent* (with the PDE)

Our Starting Point: {Diffeomorphic Mapping} = {Inverse PDE}

 $\min_{u,v,A} R(u,v,A)$

s.t. $\nabla \cdot [A(x)\nabla u(x)] = 0, \quad \forall x \in \Omega \setminus \partial \Omega$ $\nabla \cdot [A(x)\nabla v(x)] = 0, \quad \forall x \in \Omega \setminus \partial \Omega$ $u(x) = b_u(x), \quad \forall x \in \partial \Omega$ $v(x) = b_v(x), \quad \forall x \in \partial \Omega$ $n(x)^{\mathsf{T}}[A(x)\nabla u(x)] = g_u(x), \quad \forall x \in \partial \Omega$ $n(x)^{\mathsf{T}}[A(x)\nabla v(x)] = g_v(x), \quad \forall x \in \partial \Omega$ $\min_{u,v} E(u,v)$

s.t. det $[\nabla u(\mathbf{x}) \nabla v(\mathbf{x})] > 0$

(barrier methods, interior point methods)

Constraints: copy-paste previous conditions

Objective R: is some regularizer or energy E(u, v)

PDE-constrained optimization

complicated but *more efficiently* solvable with our method!

many unsuccessful attempts (augmented Lagrangian etc., too slow)

Application: Bijective Parameterizations Optimizing Different Energies

• By choosing different regularizer *R*.



Swap Two Point Landmarks Using Our Method

- The discrete solution can be quite different from the continuous one
- Require careful consideration from discrete geometry







our setting: no up-sampling

swap two point landmarks

with an up-sampled mesh

Our Framework Leads to a Family of Methods

• Different functionals in our framework provide variant means solving the problem



Experiments

- Our method is extremely robust and fast
- Pass a challenge with 11k tests [Du et al. 2021]; up to 1000 faster



diffeomorphisms by our method, shapes already cut into disk topology







Inverse PDEs v.s. Geometric Optimization



[Lipman 2012]. *E_{MIPS}=1.03.* max=3.00, mean=2.63. aliasing patterns: triangulation-sensitive



Ours (si-log). *E_{MIPS}*=1.01. max=3.11, mean=2.63. smoother 50X faster Ours (tanh+MIPS). *E_{MIPS}*=0.55. max=27.83, mean=2.00. and support many objectives

Application: Collision Avoidance by way of Injectivity

- Task: Put the Cheeseman in an hour-glass
- Ours: warp the shape + surrounding space
 - ensuring shape-mesh + air-mesh inversion-free avoids penetration / self-intersection
- Result: a unified engine for physics & collision







Why Inverse PDE is Better? Evidences from Inverse Problems of PDEs

Problem: $\min_{A,u,v} E(A, u, v)$

Solvers: operating in the space of:

prior work in (u, v):
 min { min E(A, u, v) }
 non-smooth & grad vanish

 ours in (A) min { min E(A, u, v) } a tight upper bound smooth & C[∞] differentiable



A Geometry Approach for Topological Constraints

"Fast Quasi-Harmonic Weights for Geometric Data Interpolation." Yu Wang and Justin Solomon. *ACM Transactions on Graphics (TOG) 40(4). ACM SIGGRAPH 2021*. • A tool for artist to direct animation with *sparse control*



- Skinning: drive deformation by propagating transformations at skeletons to all vertices
- Fast
- Common in computer games





- Skinning: drive deformation by propagating transformations at skeletons to all vertices
- Fast
- Common in computer games





- Skinning: drive deformation by propagating transformations at skeletons to all vertices
- Fast
- Common in computer games



- Skinning: drive deformation by propagating transformations at skeletons to all vertices
- Fast
- Common in computer games





Problem: Skinning Weight Computing

- Skinning/interpolation weights: a *partition-of-unity* that decay from 1 to 0
- $w_j(\cdot)$ is a fundamental geometry quantity. E.g., $\nabla w_j(\cdot)$ defines the *foliation*. Mathematically:

(Lagrange)
$$\forall \mathbf{c}_i : w_j(\mathbf{c}_i) = \delta_{ij}$$

(Partition of unity) $\forall \mathbf{x} : \sum_{j=1}^m w_j(\mathbf{x}) = 1$
(Nonnegativity) $\forall \mathbf{x} : w_j(\mathbf{x}) \ge 0$
(No local extrema) $\forall \mathbf{x} : w_j(\mathbf{x})$ not a local extremum $j = 1, ..., m$



A generalized B-spline, for a non-Euclidean domain Ω .

 $\min_{w} \sum_{i=1}^{m} \int_{\Omega} \|\Delta w_j(\mathbf{x})\|^2$

(Lagrange) $\forall \mathbf{c}_i : w_i(\mathbf{c}_i) = \delta_{ii}$

(Partition of unity) $\forall \mathbf{x} : \sum_{i=1}^{m} w_i(\mathbf{x}) = 1$

- used extensively beyond animation
- take hours on large example

Objective: Weight Smoothness

$$i, j = 1, ..., m$$

Constraint: Desired Properties

(Nonnegativity) $\forall \mathbf{x} : w_j(\mathbf{x}) \ge 0$ j = 1, ..., m(No local extrema) $\forall \mathbf{x} : w_j(\mathbf{x})$ not a local extremum j = 1, ..., m

[Jacobson et al. 2011, 2012]

BBW or monotonic BBW (MBBW)

s.t.

Monotonicity Constraint

- Monotonicity: no local extremum away from control handles
 - A topological constraint [Jacobson et al. 2012]
 - A necessary condition for *diffeomorphic* shape morphing
- Without monotonicity ightarrow noticeable artifacts



Our Model: Quasi-Harmonic Weights

• Consider solutions to the anisotropic Laplace equations

$$\begin{split} \Delta^{\mathbf{A}} w_j(\mathbf{x}) &= 0 & \forall j = 1, ..., m \\ w_j(\mathbf{c}_i) &= \delta_{ij} & \forall i, j = 1, ..., m \\ \nabla^{\mathbf{A}}_{\mathbf{n}} w_j(\mathbf{x}) &= 0 & \forall \mathbf{x} \in \partial \Omega / \mathbf{c}, \forall j = 1, ..., m. \end{split}$$

• For any A(x), the weights $w_j(x)$ satisfy:

(Lagrange)
$$\forall \mathbf{c}_i : w_j(\mathbf{c}_i) = \delta_{ij}$$

(Partition of unity) $\forall \mathbf{x} : \sum_{j=1}^m w_j(\mathbf{x}) = 1$
(Nonnegativity) $\forall \mathbf{x} : w_j(\mathbf{x}) \ge 0$
(No local extrema) $\forall \mathbf{x} : w_j(\mathbf{x})$ not a local extremum $j = 1, ..., m$

For any A(x), the generated w(x) automatically satisfy all conditions

 $\Delta^{\mathbf{A}} = \nabla \cdot [\mathbf{A}(\mathbf{x})\nabla]$

Our Method

• Search within the family of quasi-harmonic weights $\Delta^{\mathbf{A}} = \nabla \cdot [\mathbf{A}(\mathbf{x})\nabla]$ $\min_{w} \quad \sum_{j=1}^{m} \int_{\Omega} \|\Delta w_j(\mathbf{x})\|^2$ $\min_{\mathbf{A}} \quad \sum_{j=1}^{m} \int_{\Omega} \|\Delta w_{j}(\mathbf{x})\|^{2}$ s.t. $w_{j}(\mathbf{c}_{i}) = \delta_{ij} \qquad \forall i, j = 1, \dots, m$ $\Delta^{A} w_{j}(\mathbf{x}) = 0 \qquad \forall \mathbf{x} \in \Omega, \forall j = 1, \dots, m$ $\nabla_{\mathbf{n}}^{A} w_{j}(\mathbf{x}) = 0 \qquad \forall \mathbf{x} \in \partial\Omega, \forall j = 1, \dots, m$ i, j = 1, ..., m $\forall \mathbf{c}_i : w_j(\mathbf{c}_i) = \delta_{ij}$ $\forall \mathbf{x} : \sum_{j=1}^{m} w_j(\mathbf{x}) = 1$ $\mathbf{A}(\mathbf{x}) \geq 0$ $\forall \mathbf{x} \in \Omega$ $\forall \mathbf{x} : w_j(\mathbf{x}) \geq 0$ j=1,...,m $\forall \mathbf{x} : w_j(\mathbf{x}) \text{ not a local extremum} \quad j = 1, ..., m$ **PDE-constrained optimization** • not necessarily easier • but we find an efficient solver

Evaluation: Timing

Example			Smoothness Energy					Time (Sec.)		
Mesh	#Hdl.	#Ele.	BBWA	MBBW	Ours, $k = 20$	k = 10	BBWA	MBBW	Ours $(k=10)$	num. fact.
Beast	15	443442	1	1.001	0.997	1.016	707.8	840.6	12.6	0.16
Bunny	14	531392	1	0.997	1.001	1.026	2430.6	11111.3	18.0	0.41
Raptor	15	367966	1	1.002	0.992	1.051	640.1	720.3	14.9	0.11
Elephant	17	516858	1	0.996	0.987	1.008	1092.3	2379.5	14.4	0.16
Dragon	17	1187670	1	*	0.992	1.023	5022.1	*	54.8	0.60
Image	27	6952	1	0.995	0.984	1.005	7.78	17.70	0.18	0.0009
Brick	2	78529	1	1.000	0.981	1.084	17.8	53.7	1.21	0.03
Tibiman	16	84125	1	0.991	0.986	0.995	91.9	274.1	1.80	0.025

• Our inverse PDE solver: orders-of-magnitude faster than previous methods

BBW [Jacobson et al. 2011] MBBW [Jacobson et al. 2012]



BBWA (0.68 hr)

MBBW (3.1 hr)

Ours (0.3 min)

57

Geometric Computing beyond the Laplacian

- Optimization in a larger space of Laplacians.
- Design Laplacian-like operators.
- Learn Laplacian-like operators from data.

Shape Classification backed by Modern Geometric Analysis

"Steklov Spectral Geometry for Extrinsic Shape Analysis" Yu Wang, Mirela Ben-Chen, Iosif Polterovich, Justin Solomon ACM Transactions on Graphics 38(1)

Parallel Roles of Geometry and Image Processing

 $= C_1$ 1

- Image processing & analysis
 - input: 2D array
- Image classification

- Geometry processing & analysis
 - input: 2-dim manifold in 3D
- Shape classification

• Local features: SIFT etc.



• Local features: curvatures etc.



[Lombaert et al. 2013] 60

Geometric Computing Tasks

- Task: shape analysis and 3D vision
 - shape classification



- Approach
 - Theory: insights taken from modern **spectral geometry**
 - Tools: borrowed from **computational electromagnetics**

Mathematically Justified Embedding: Laplacian Eigenvalues

• Solve the eigenvalue problem

$$\Delta \phi_i = \lambda \phi_i$$



• $(\lambda_1, \lambda_2, \dots, \lambda_k) \in \mathbb{R}^k$: Shape2Vector scheme, mathematically justified

Intrinsic/Laplacian approaches are invariant to isometry ("pose invariant") Real-world objects are usually subject to (near-) isometries



Laplacian Shape Classification



(Intrinsic) Laplacian can be Counterintuitive



• Same Laplacian operator \rightarrow same Laplacian eigenvalues

(Intrinsic) Laplacian Information is Incomplete



Intrinsic geometry: any origami is equivalent to a piece of flat paper!

Laplacian Lacks Robustness





- Should be identical
- But completely different are their Laplacian eigenvalues---Why?

Laplacian is Sensitive to Topological Noises





• The surface connected to legs

- Not connected. The surface only *touches* the legs
- Topology errors → completely different Laplacian eigenvalues

Our Solution: Replace the Laplacian Δ with Operator S

$\begin{array}{c} \Delta \rightarrow \mathcal{S} \\ L \rightarrow \mathcal{S} \end{array}$

Our Solution: DtN Operator and Steklov Eigenvalue Problem

• Discrete Dirichlet-to-Neumann (DtN) operator: $S \in \mathbb{R}^{n \times n}$

n: number of vertices



Definition

 \mathcal{S} : the Dirichlet-to-Neumann operator. The Steklov eigenvalue problem

 $\mathcal{S}\psi = \lambda\psi$



Theorem (Lassas 2001)

Denote $\Omega_1, \Omega_2 \subseteq \mathbb{R}^3$ as two domains, and $\alpha : \Omega_1 \to \Omega_2$ is a bijection. Under proper assumptions, if the two domains have the same Dirichlet-to-Neumann operators (under map composition), then α must be a rigid motion.

For smooth domains in \mathbb{R}^3 , the Steklov heat kernel admits the asymptotic expansion [Polterovich and Sher 2015]

$$e^{-t\mathcal{S}}(x,x) = \sum_{i=0}^{\infty} e^{-t\lambda_i} \phi_i(x)^2 \sim \sum_{k=0}^{\infty} a_k(x) t^{k-2} + \sum_{l=1}^{\infty} b_l(x) t^l \log t,$$

H(x): mean curvature K(x): Gaussian curvature

$$a_0(x) \equiv \frac{1}{2\pi}$$

$$a_1(x) = \frac{H(x)}{4\pi}$$

$$a_2(x) = \frac{1}{16\pi} \left(H(x)^2 + \frac{K(x)}{3} \right)$$




Spectral clustering with geodesics / heat kernels \rightarrow shape segmentation

Level sets of Steklov eigenfunctions conform to mean curvatures

Laplacian Segmentation: Much Worse than Ours



Heat Kernel Signature $h_t(x)$



Robust to Noises



76

Boundary Approach: Robust to Non-watertight Surface

Steklov eigenfunctions: stable to topological changes for open surfaces





Laplacian Steklov Shape Classification



Summary: Mathematical Geometric Processing: Laplacian & Beyond

- A Tutorial on Laplacian
- Quasi-harmonic maps
 - search for a deformed Laplacian
- Extrinsic shape analysis
 - design a new operator
- Learn operator kernel from data
 - for a different operator kernel
- Future: potentially every tasks in geometric computing...
- Links to the paper: <u>https://wangyu9.github.io/</u>
- For further questions: wangyu9@mit.edu





