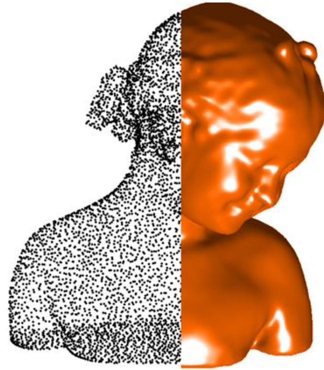
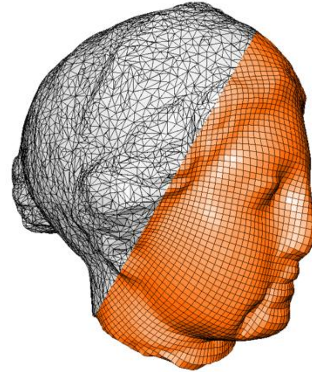
Digital Geometry Processing

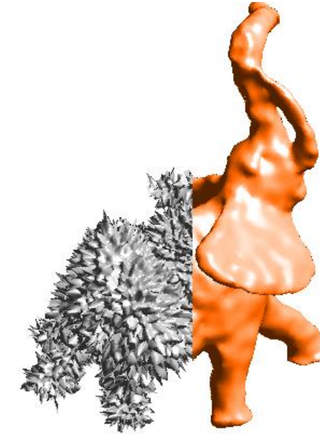
Overview of Geometry Processing Tasks



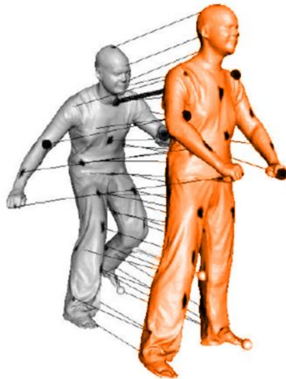
[reconstruction]



[remeshing]

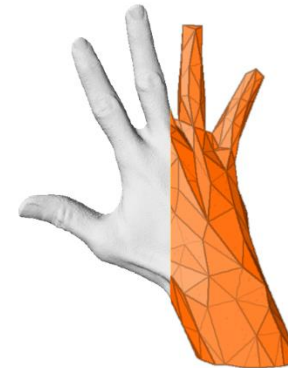


[filtering]



[shape analysis]

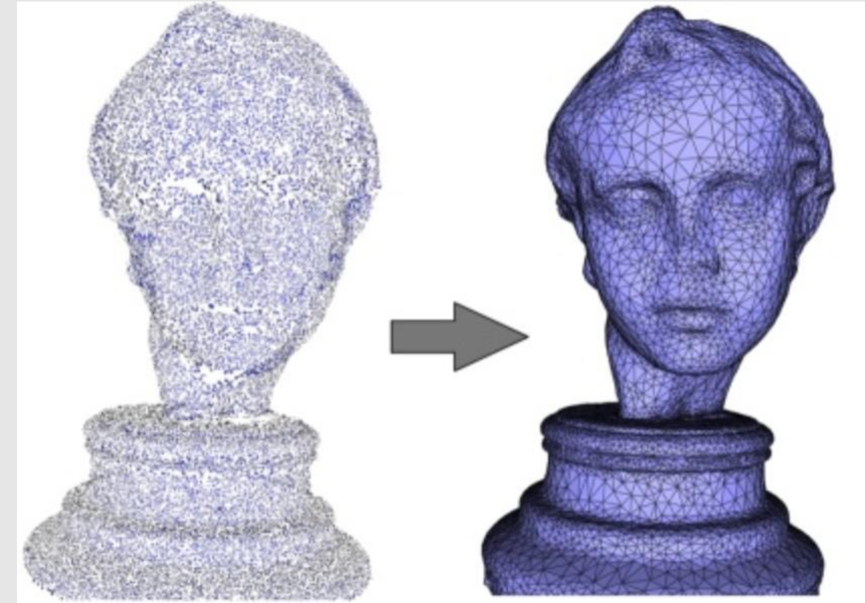
[parameterization]

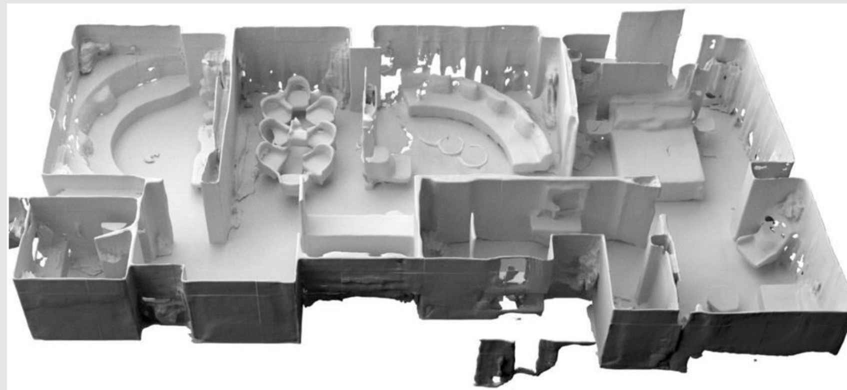


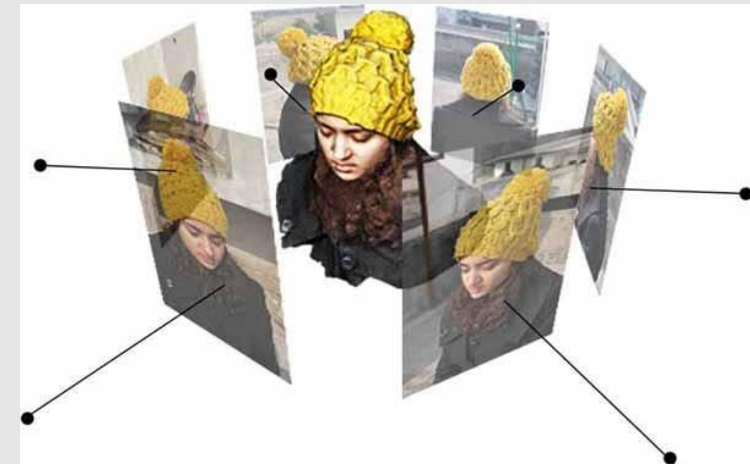
[compression]

Geometry Processing: Reconstruction

- Given samples of geometry, reconstruct surface
- **Data:** What are "samples"?
  - Points & normals
  - Image pairs / sets (multi-view stereo)
  - Line density integrals (MRI/CT scans)
- **Algorithm:** How do you get a surface?
  - Silhouette-based (visual hull)
  - Voronoi-based (e.g., power crust)
  - PDE-based (e.g., Poisson reconstruction)
  - iso-surfacing (marching cubes)

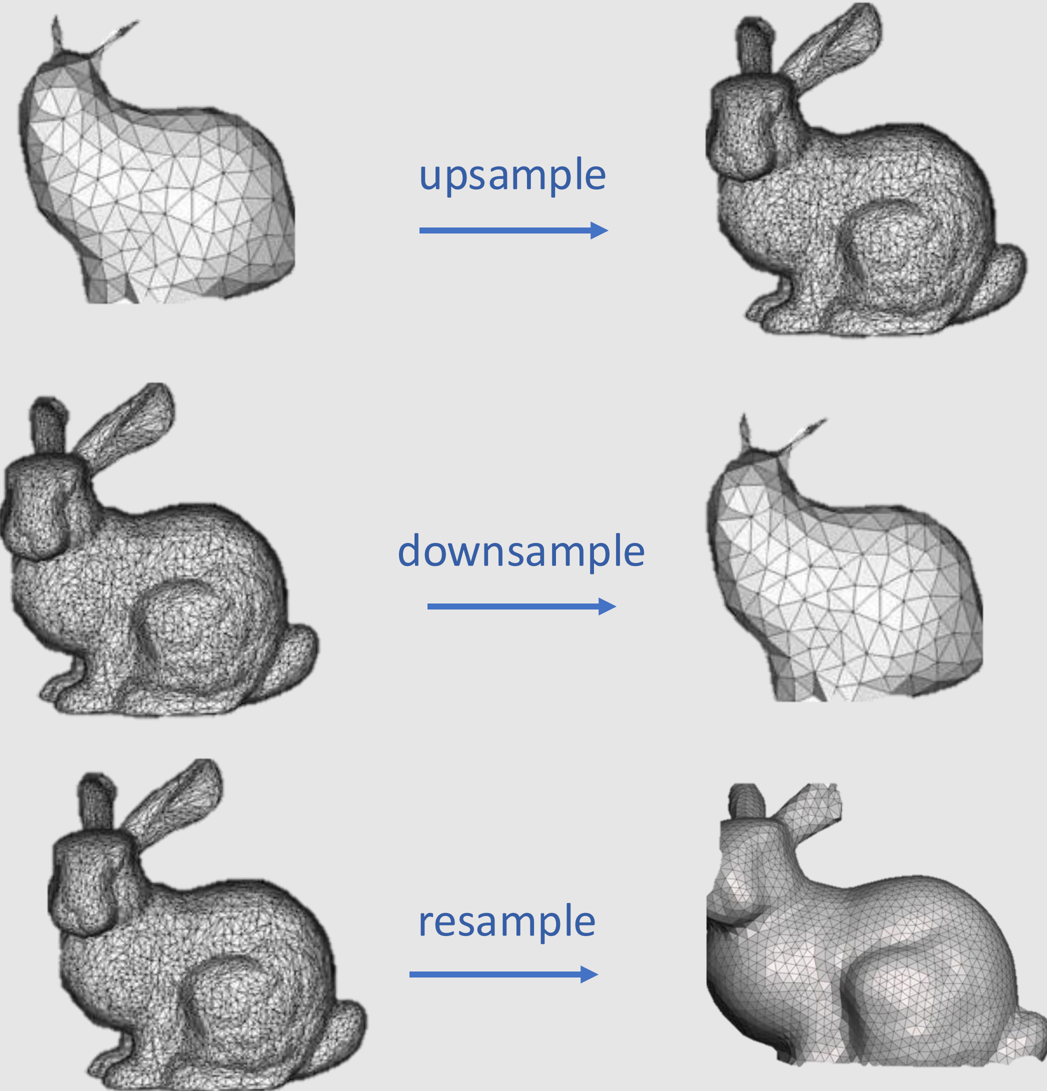






Geometry Processing: Remeshing

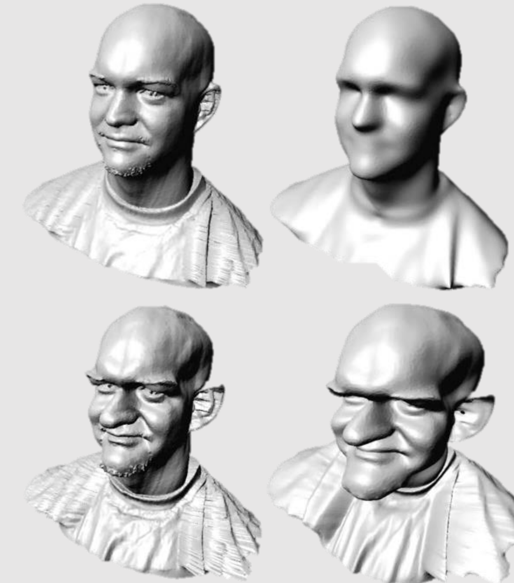
- **Upsampling**: increase resolution via interpolation, e.g.
  - Subdivision
- **Downsampling**: decrease resolution via averaging, e.g.
  - quadric error metric-based iterative decimation
  - Clustering-based mesh coarsening
- **Resampling:** modify sample distribution to improve quality, e.g.
  - Isotropic remeshing
  - Convert triangle mesh to quad mesh

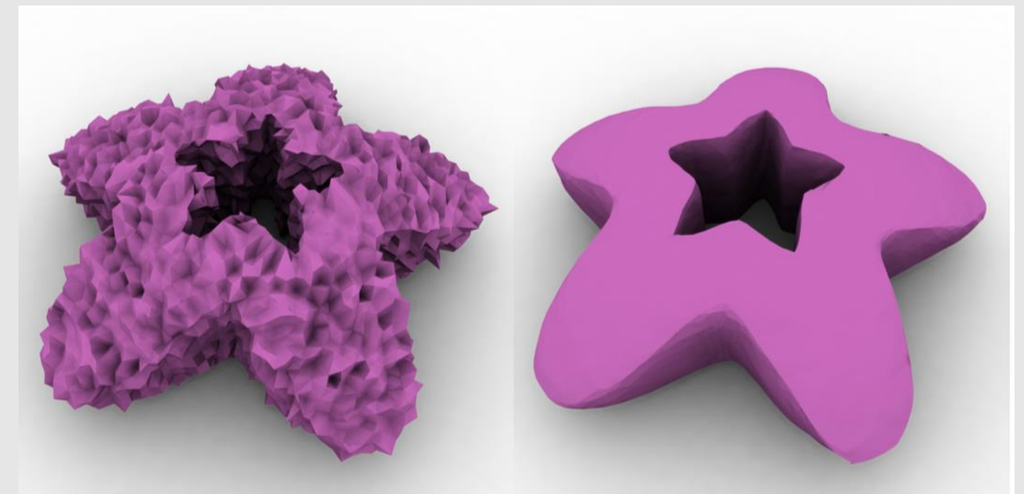


Geometry Processing: Filtering

- Remove noise, or emphasize important features (e.g., edges)
  - Curvature flow
  - Bilateral filtering
  - Spectral filtering

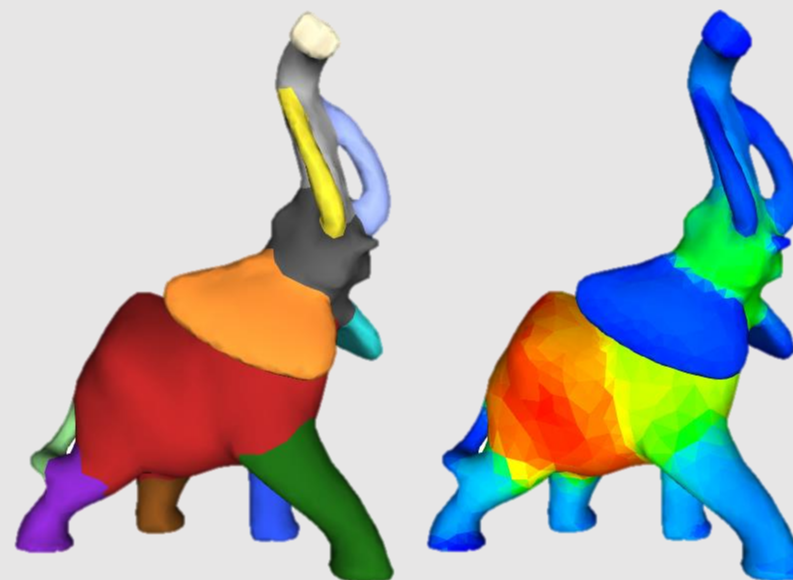
- Useful for cleaning up noisy 3D scans
  - **Example:** Kinect
    - Search for key facial components while smoothening out artifacts in between

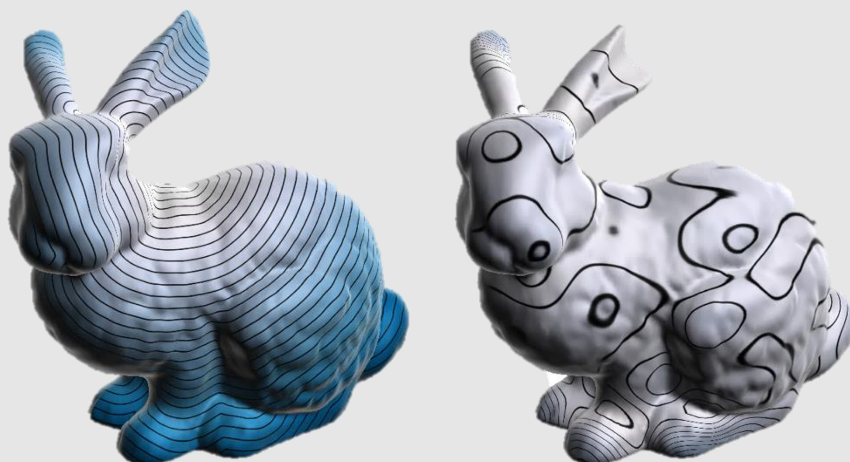


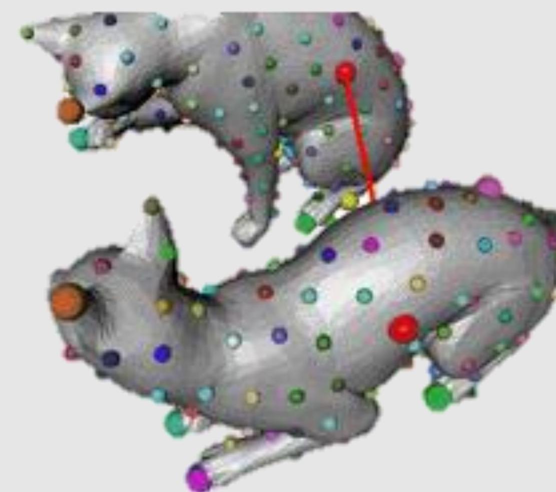


Geometry Processing: Shape Analysis

- Identify/understand important semantic features
  - Segmentation
  - Correspondence
  - Symmetry detection
  - Alignment
    - **Objective:** Compute similarities between two meshes
- Starting to become AI-driven





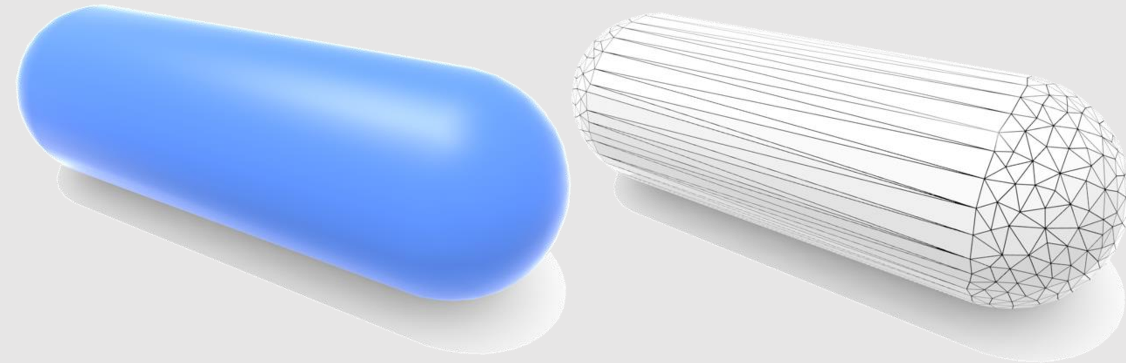


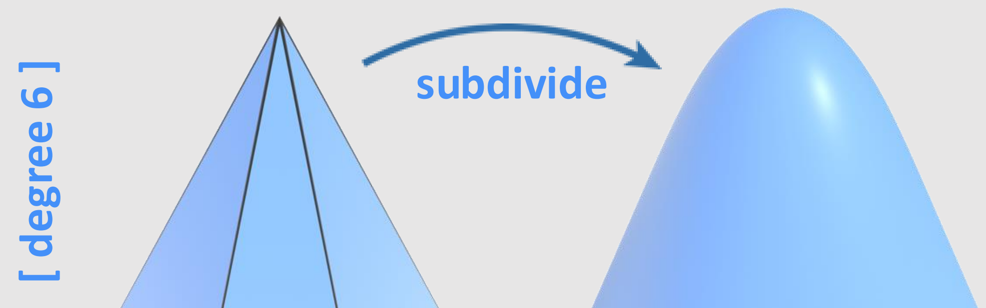
But what makes a good mesh?

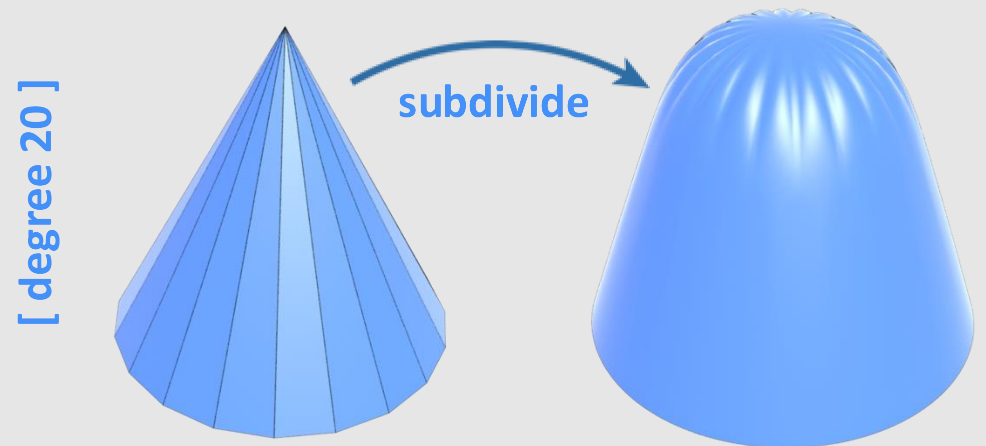
- Good Geometry
- Geometric Subdivision
- Geometric Simplification
- Geometric Remeshing
- Geometric Queries

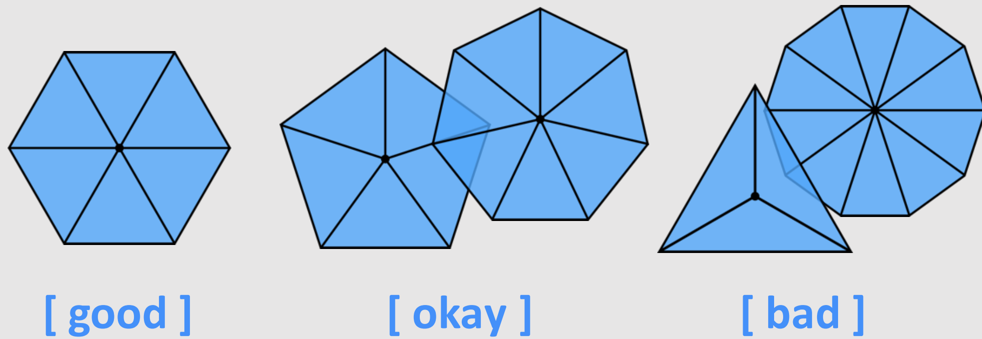
A Good Mesh Has...

- **Good approximation of original shape**
  - Keep elements that contribute shape info
  - More elements where curvature is high
- **Regular vertex degree**
  - Degree 6 for triangle mesh, 4 for quad mesh
    - Better polygon shape
    - More regular computation (less numerical issues)
    - Smoother subdivision





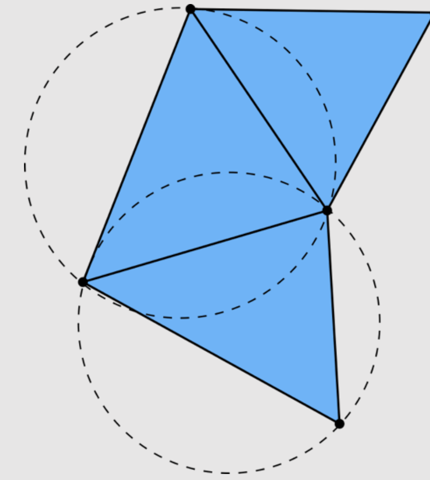




A Good Mesh Has…

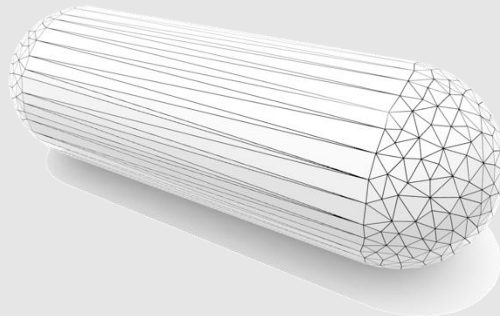
- **Good triangle shape**
  - All angles close to 60 degrees

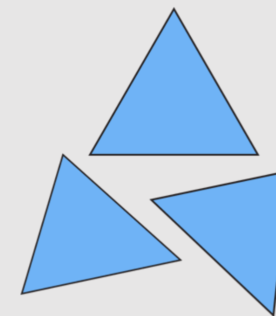
  - More sophisticated condition: **Delaunay**
    - For every triangle, the unique circumcircle (circle passing through all vertices of the triangle) does not encase any other vertices
    - Many nice properties:
      - Maximizes minimum angle
      - Smoothest interpolation



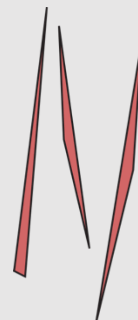
[delaunay]

- **Tradeoff:** sometimes a mesh can be approximated best with long & skinny triangles
  - Doesn't make the mesh Delaunay anymore
  - **Example:** cylinder





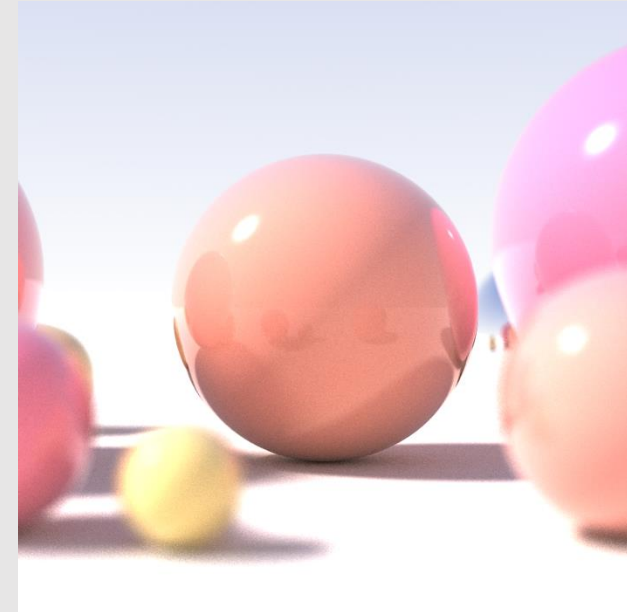
[good]

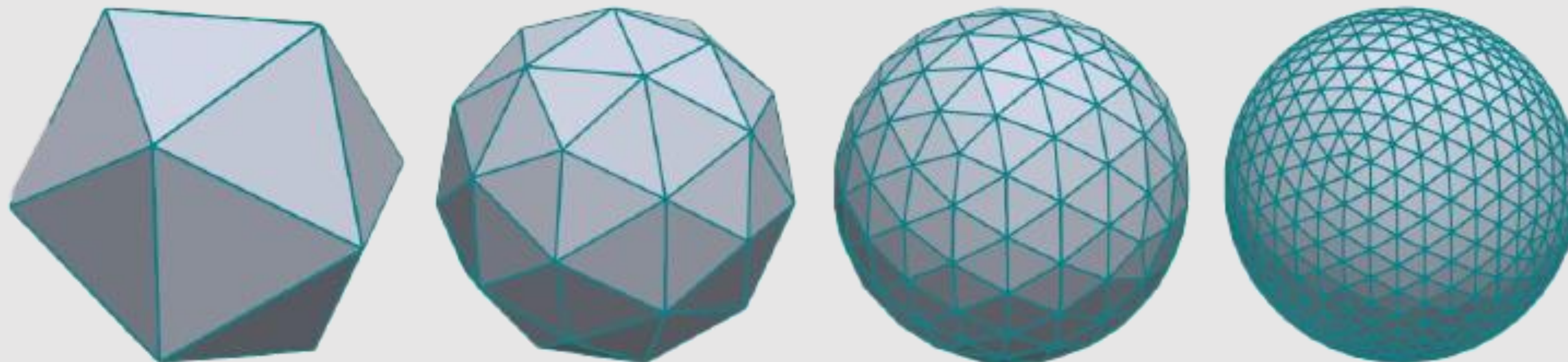


[bad]

A Good Mesh Has…

- **Good approximation of original shape (revisit)**
  - Placing vertices on a sphere and linearly interpolating is not enough
    - Adding more vertices yields better approximation, but now too much data to store/process!
  - We can better approximate the **surface normal**, so that the appearance is smoother despite coarse geometry





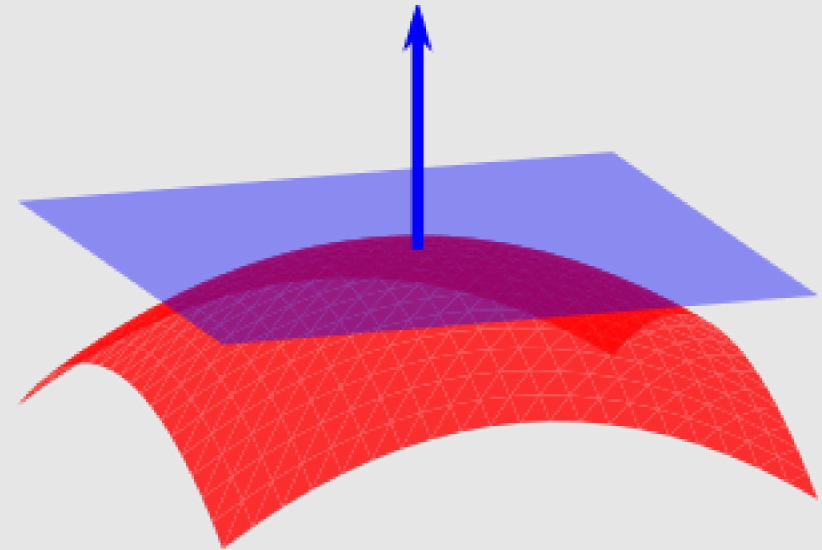
Surface Normals

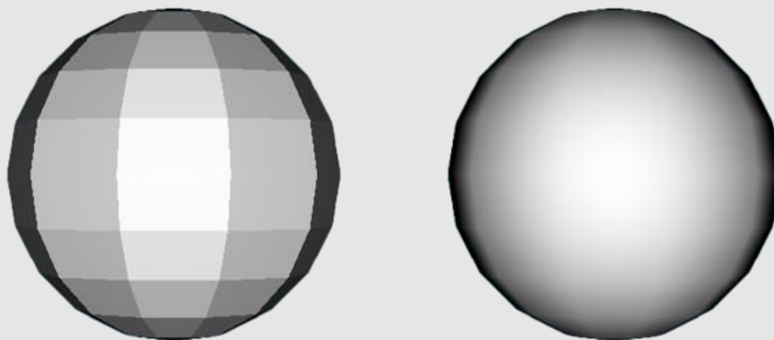
- A **surface normal** is a vector that is perpendicular to the surface at a given point
  - The surface normal for a surface $z = f(x, y)$ at point (x', y') is:

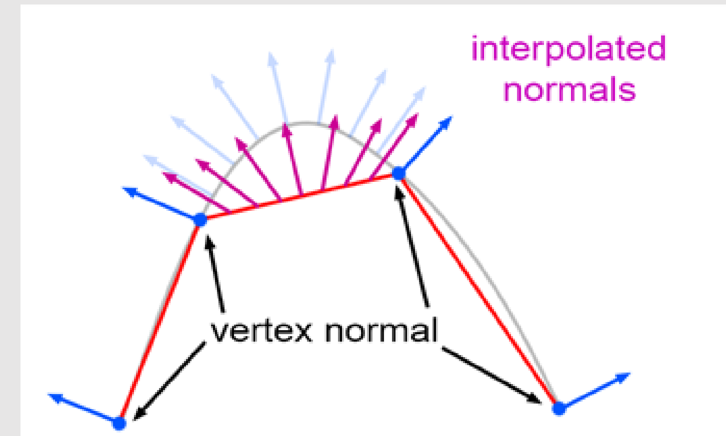
$$N_s = \begin{bmatrix} f_x(x', y') \\ f_y(x', y') \\ -1 \end{bmatrix}$$

  - Value assigned per-vertex

- Surface normal are interpolated via-barycentric coordinates and normalized to provide the appearance of curvature during rendering



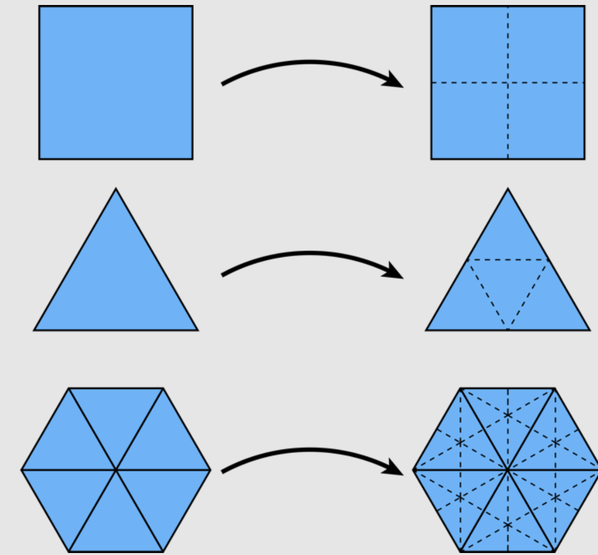


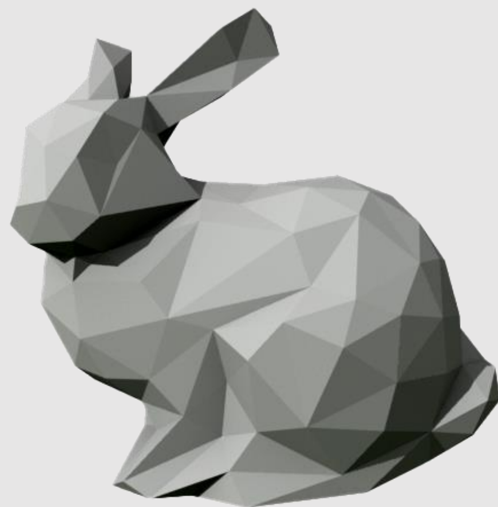


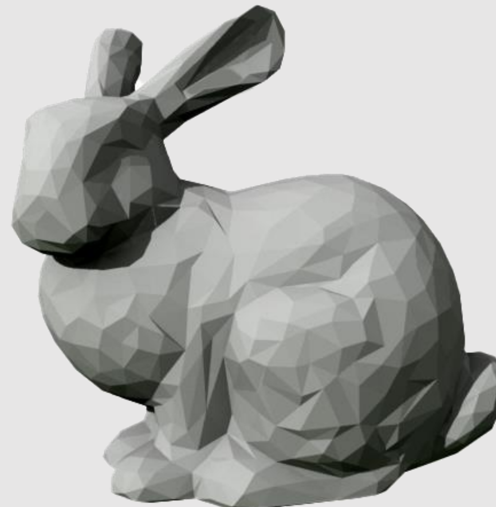
- ~~Good Geometry~~
- Geometric Subdivision
- Geometric Simplification
- Geometric Remeshing
- Geometric Queries

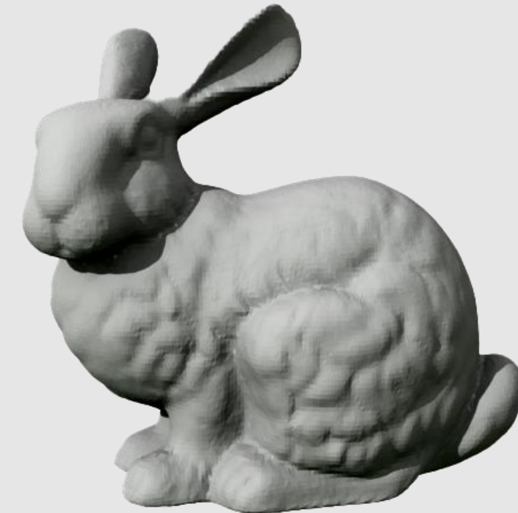
Subdivision

- Subdivision is the process of **upsampling** a mesh
- General formula:
  - **Split Step:** split faces into smaller faces
  - **Move Step:** replace vertex positions/properties with weighted average of neighbors



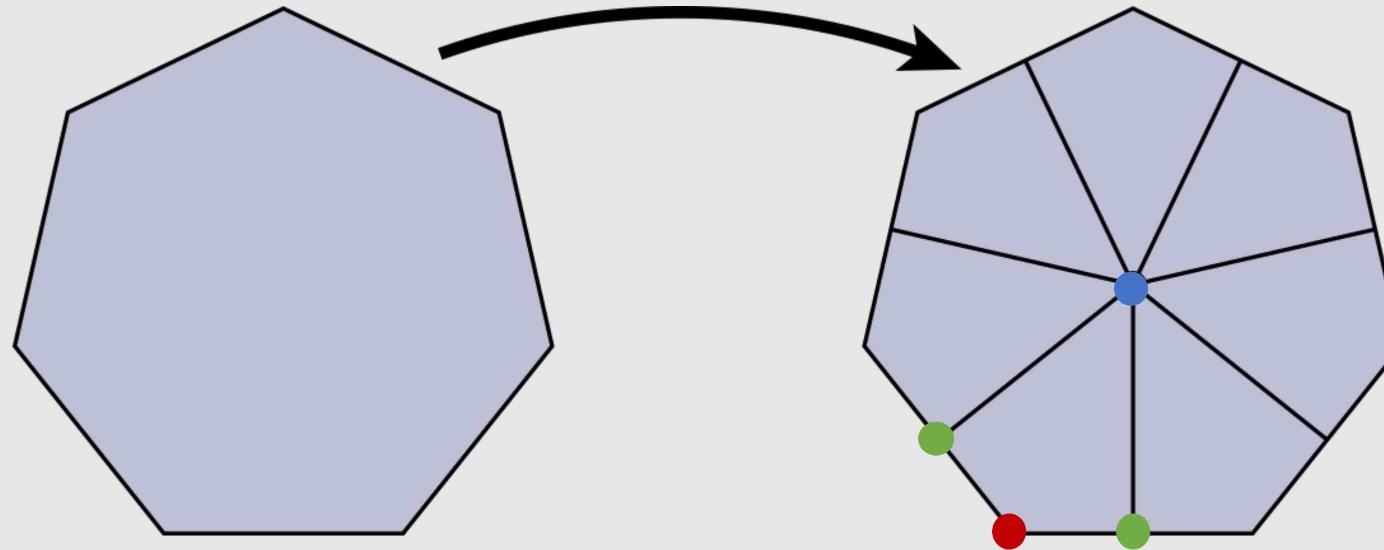






Linear Subdivision [Split Step]

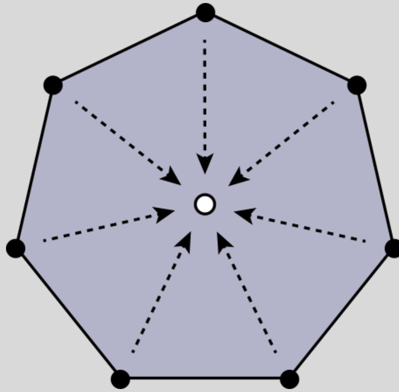
- Split every polygon (any # of sides) into quadrilaterals



- Each new quadrilateral now has:
  - [face coords] : 1 new vertex from the mesh face center
  - [edge coords] : 2 new vertices from the mesh edges
  - [vertex coords] : 1 new vertex from the original mesh face

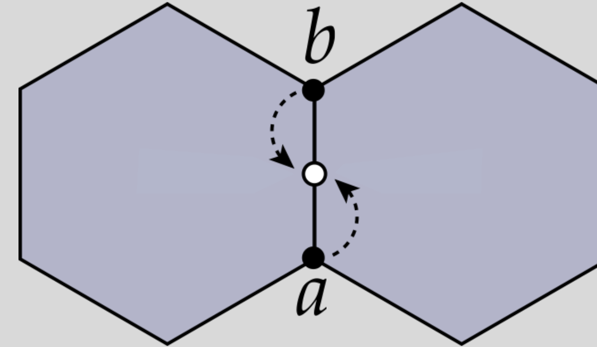
Linear Subdivision [Move Step]

Step 1: Face Coords



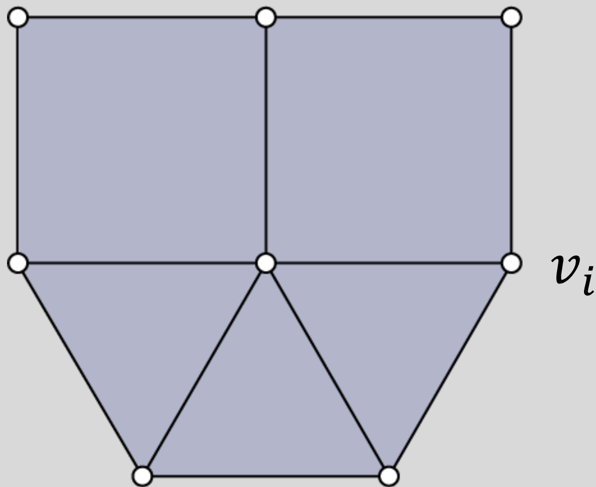
$$\frac{1}{n}\sum_i p_i$$

Step 2: Edge Coords



$$(a + b) / 2$$

Step 3: Vertex Coords



v_i

$$v_i = v_i$$

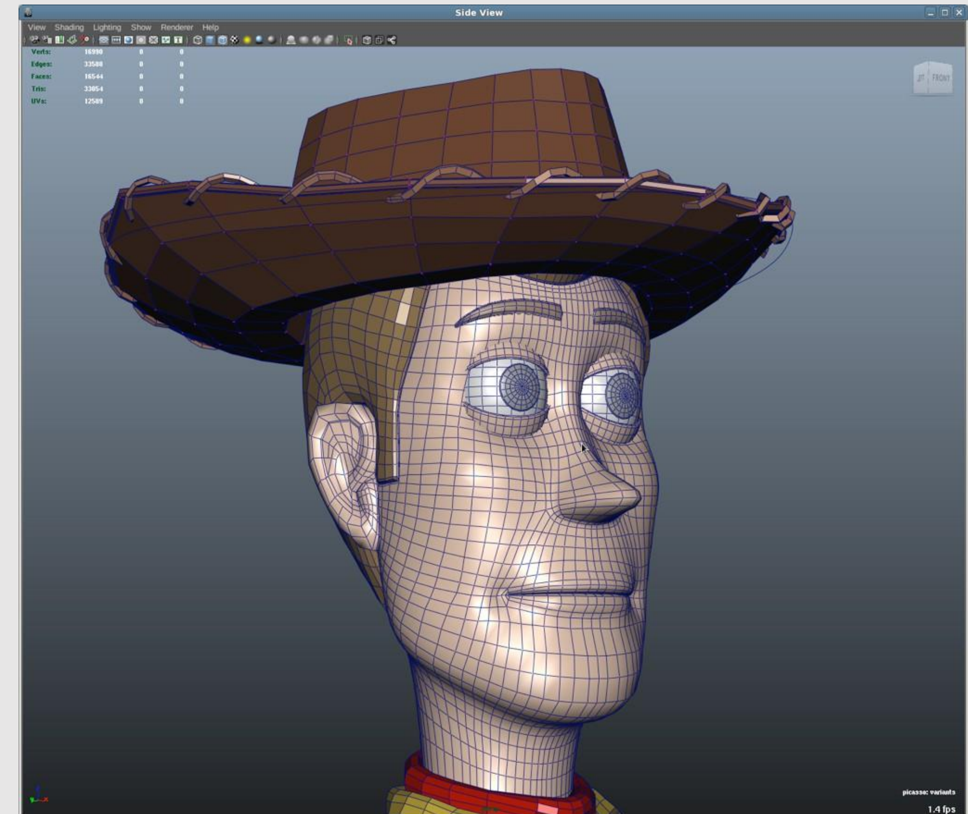
But linear subdivision does not change the mesh geometry much.

-- Almost all added vertices, edges, and faces lie exactly on the mesh surface.

Catmull-Clark Subdivision

- In 1978, Edwin Catmull (Pixar co-founder) and Jim Clark wanted to create a generalization of **uniform bi-cubic b-splines** for 3D meshes
  - We will cover what this means in a future lecture :)

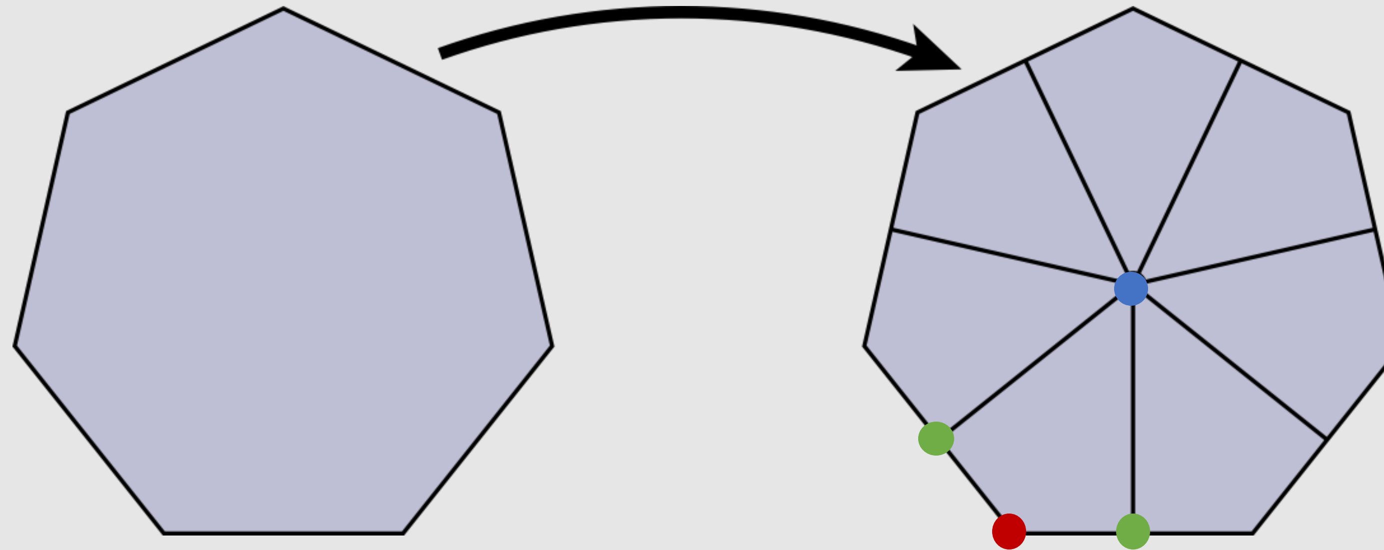
- Became ubiquitous in graphics
  - Helped Catmull win an Academy Award for Technical Achievement in 2005



OpenSubdiv V2 (2018) Pixar

Catmull-Clark Subdivision [Split Step]

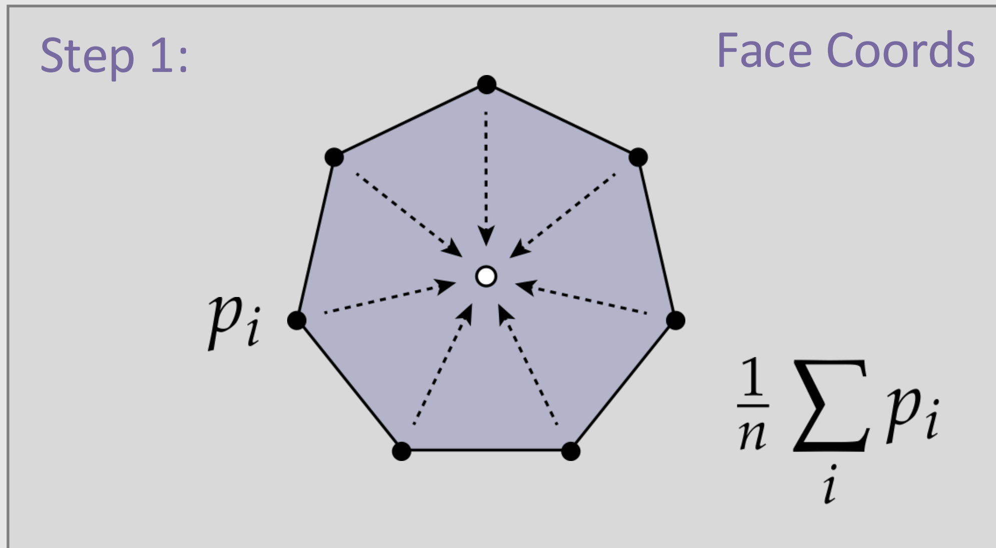
- Split every polygon (any # of sides) into quadrilaterals

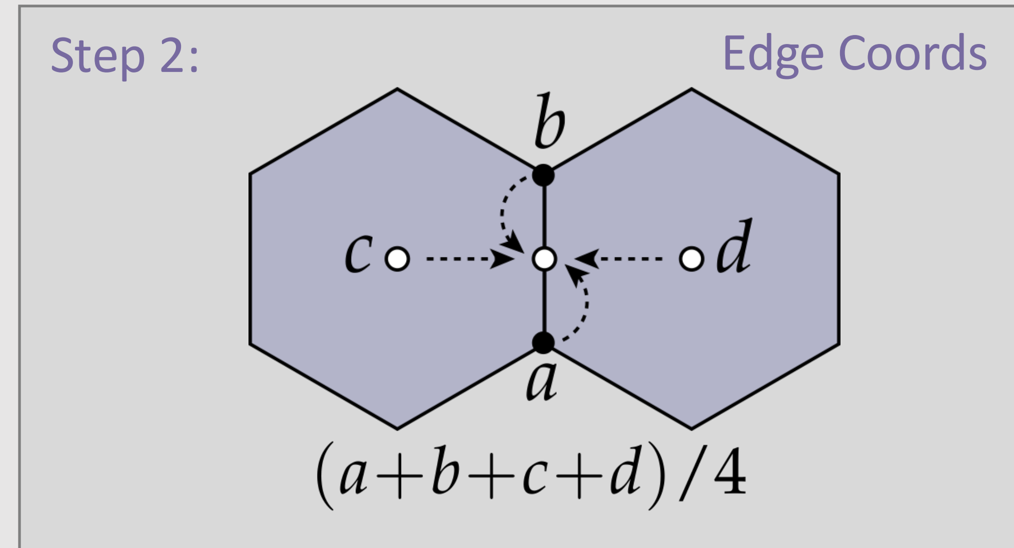


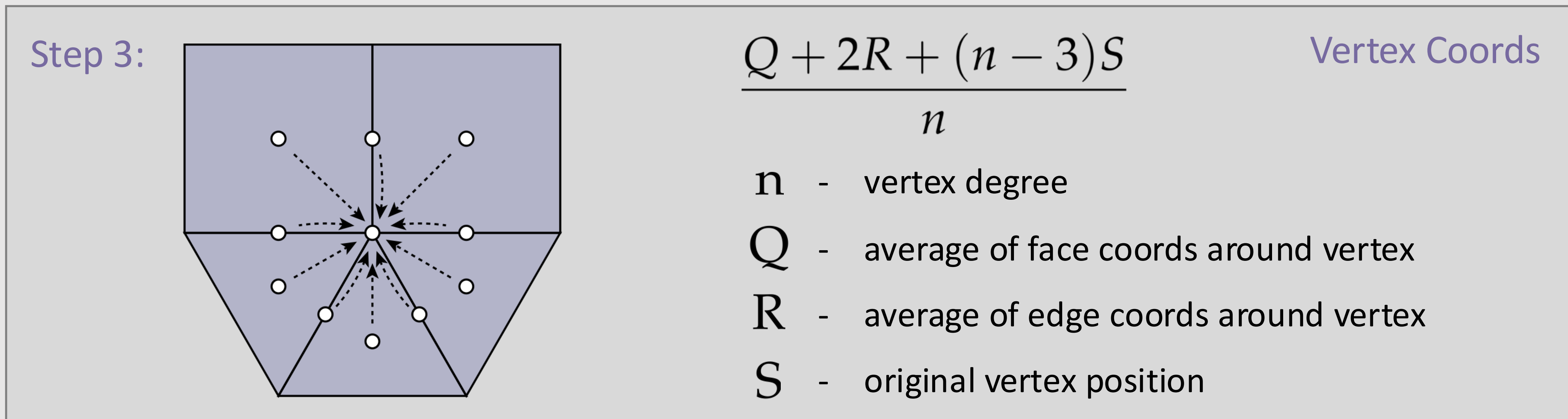
- Each new quadrilateral now has:
  - [face coords] : 1 new vertex from the mesh face center
  - [edge coords] : 2 new vertices from the mesh edges
  - [vertex coords] : 1 new vertex from the original mesh face

No different than Linear Subdivision!

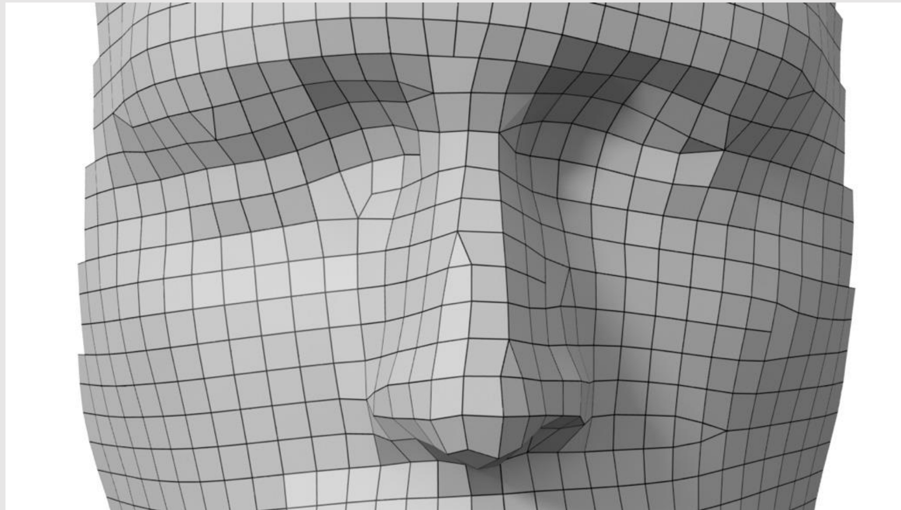
Catmull-Clark Subdivision [Move Step]

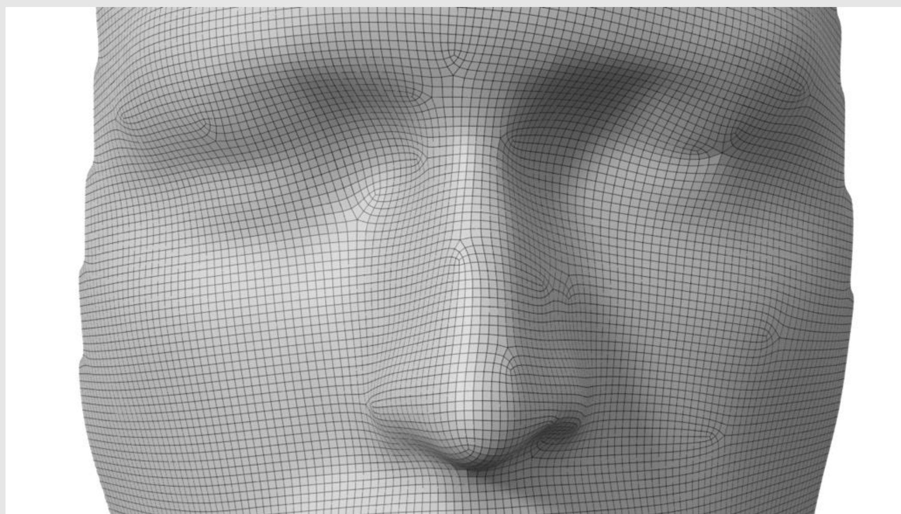






Catmull-Clark Subdivision [Quads]



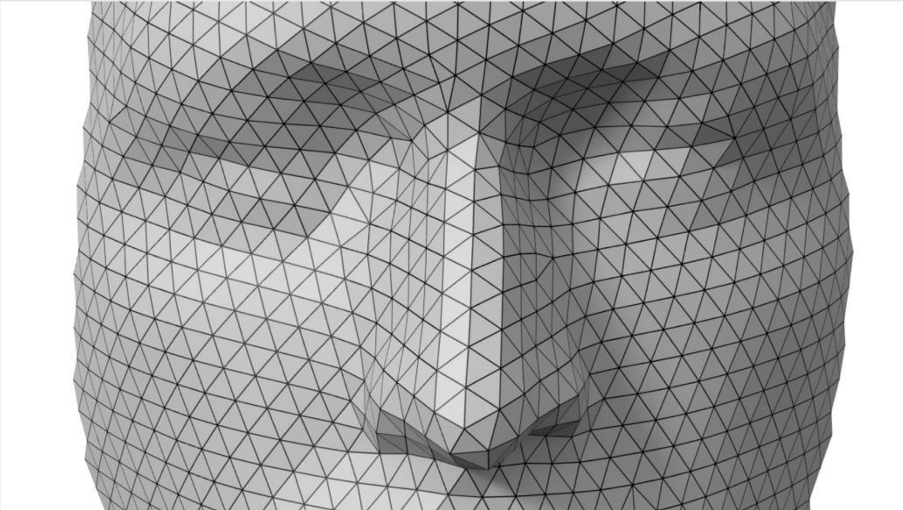


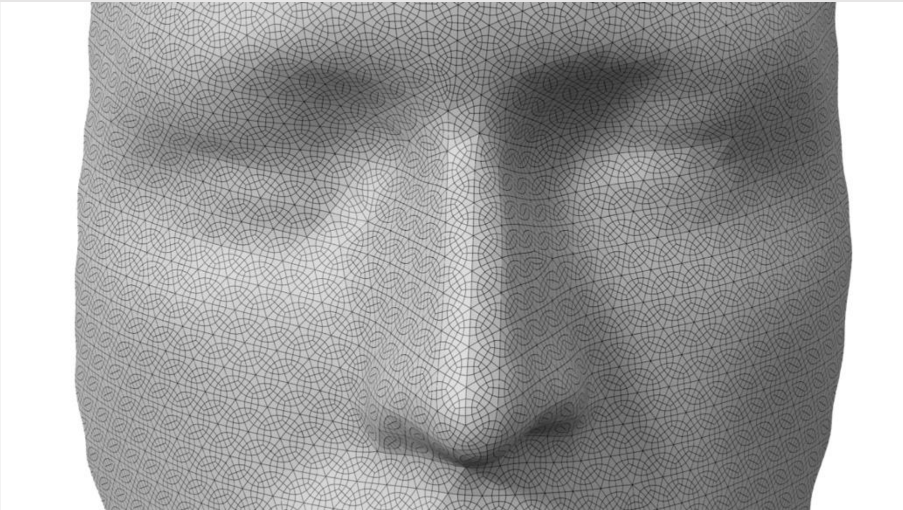
Few irregular vertices

Smoothly-varying surface normals

Smooth reflections/caustics

Catmull-Clark Subdivision [Triangles]





Many irregular vertices

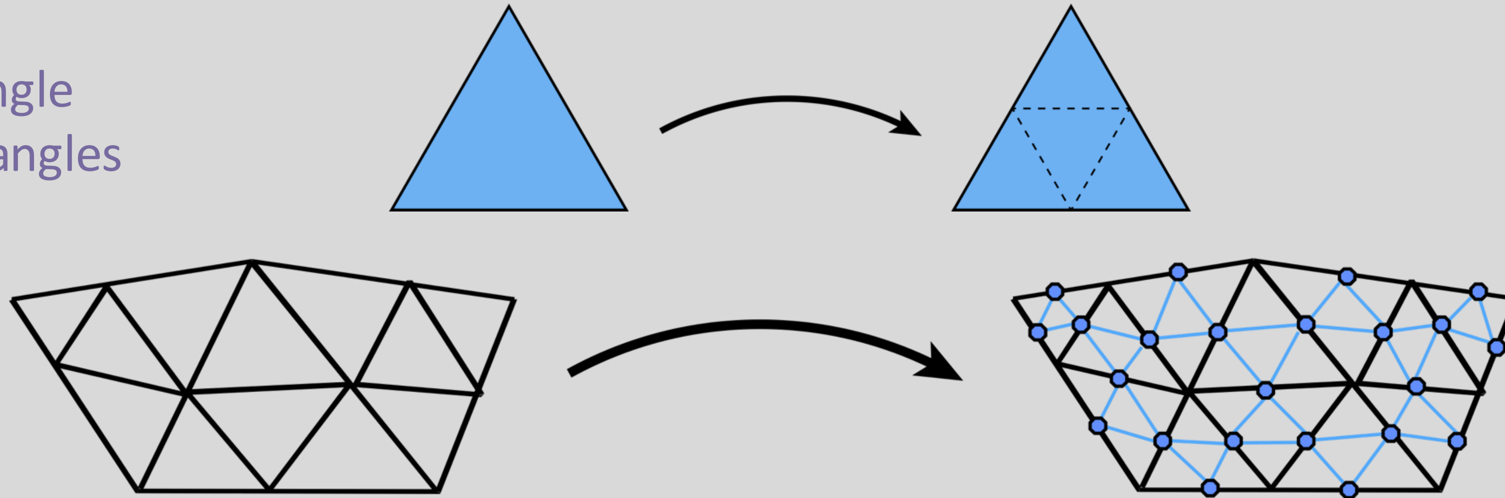
Erratic surface normals

Jagged reflections/caustics

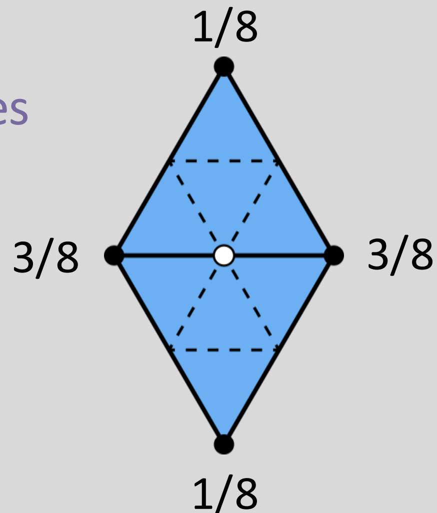
Is there a better subdivision scheme we can use for triangulated meshes?

Loop Subdivision

Step 1:
Split triangle
into 4 triangles

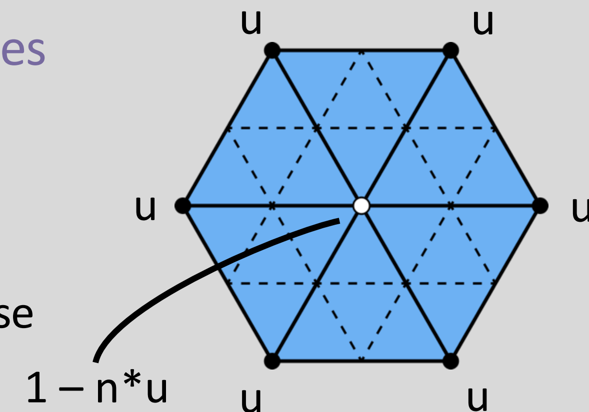


Step 2:
Assign coordinates
for new vertices



Step 3:
Assign coordinates
for old vertices

n - vertex degree
u - 3/16 if n=3
3/(8n) otherwise



Loop Subdivision

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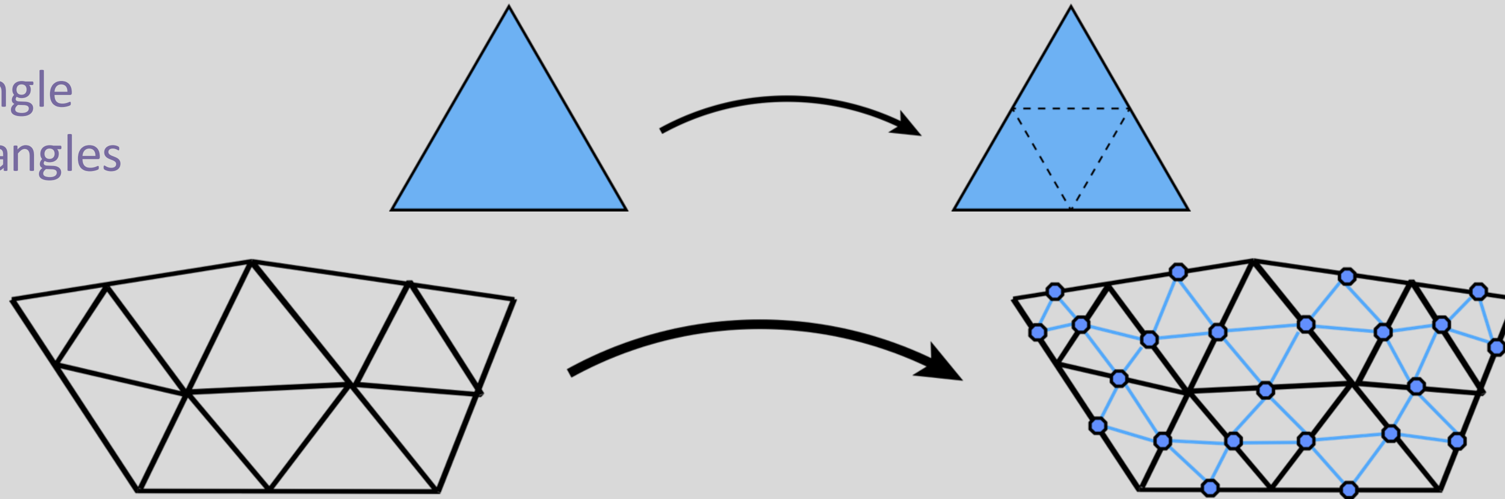
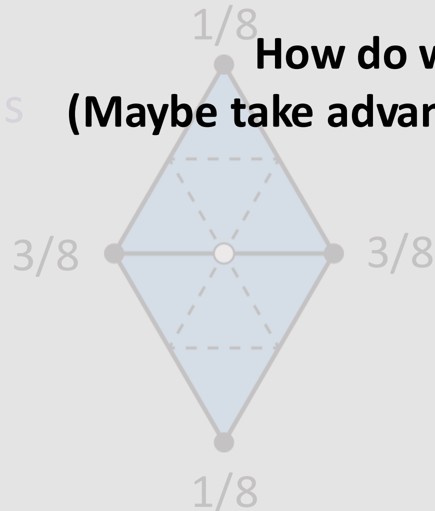


Step 2:
Assign new coords



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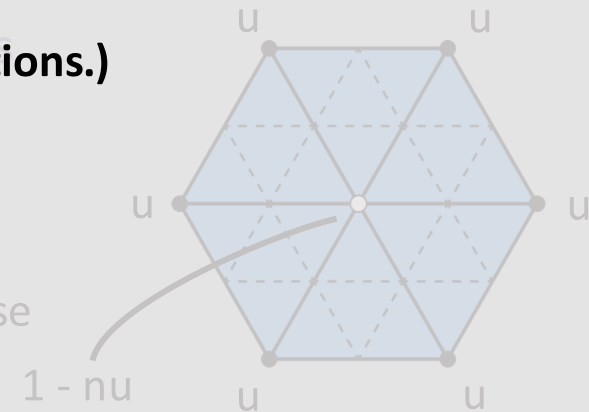
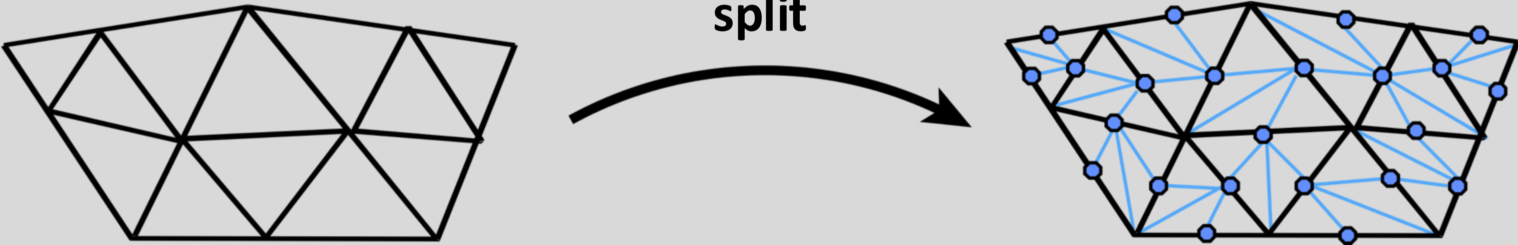


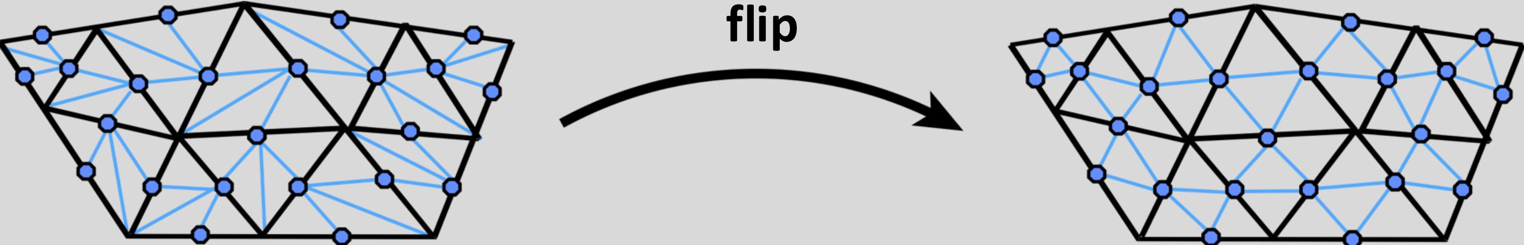
How do we conveniently do Step 1?
(Maybe take advantage of local geometric operations.)

Loop Subdivision Using Local Ops

Step 1:
Split all edges in any order

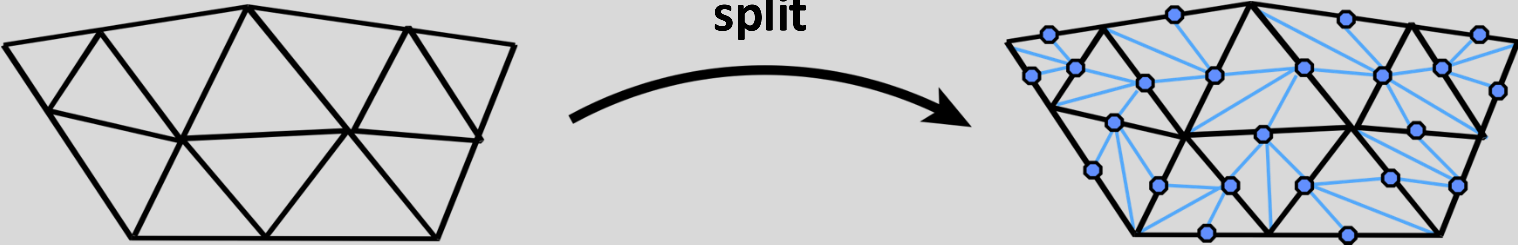


Step 2:
Flip new edges until they touch two new vertices



Loop Subdivision Using Local Ops

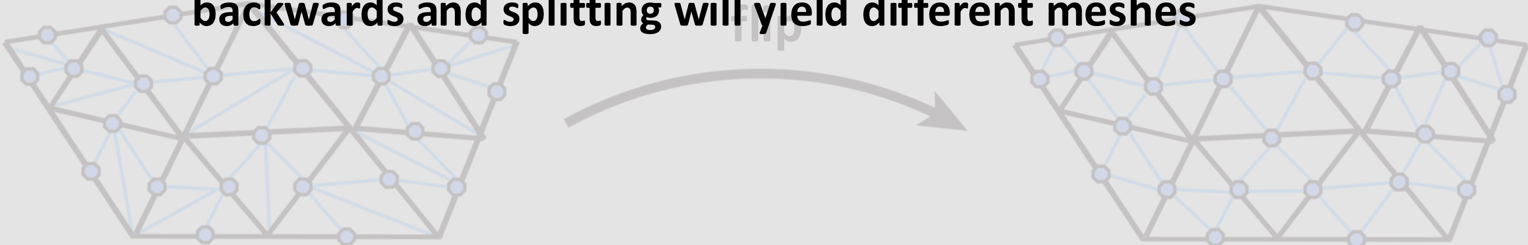
Step 1:
Split all edges in any order



Step 2:
Flip new edges until they touch two new vertices

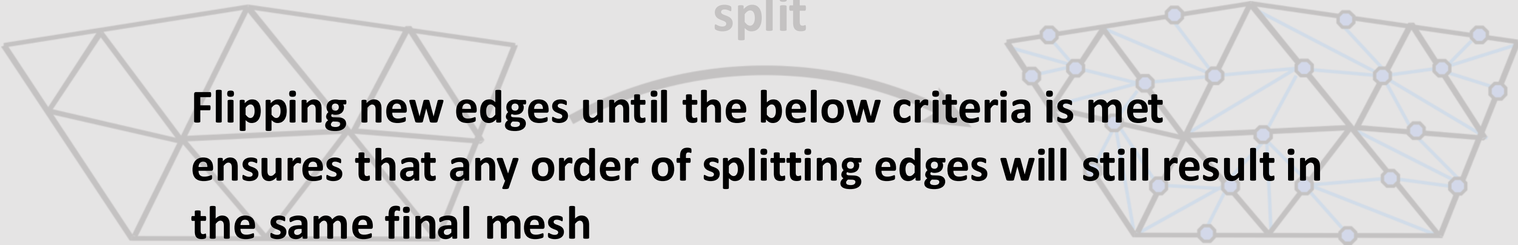
The order we traverse the edges and split them matter!

Traversing edges forward and splitting vs traversing them backwards and splitting will yield different meshes



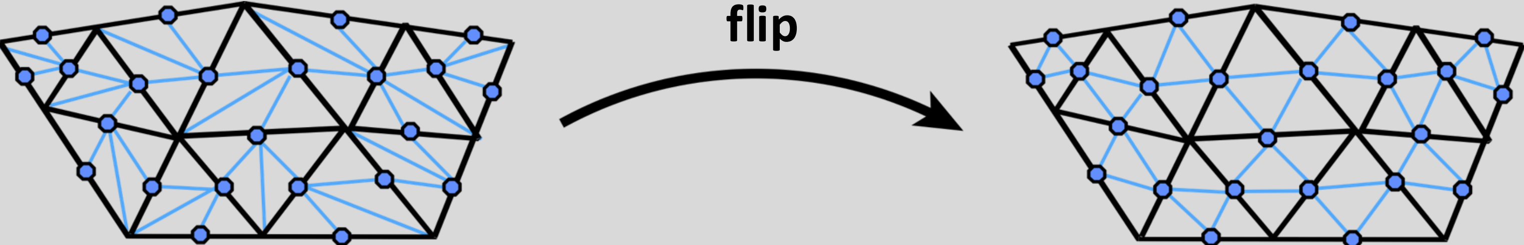
Loop Subdivision Using Local Ops

Step 1:
Split all edges in any order



Flipping new edges until the below criteria is met ensures that any order of splitting edges will still result in the same final mesh

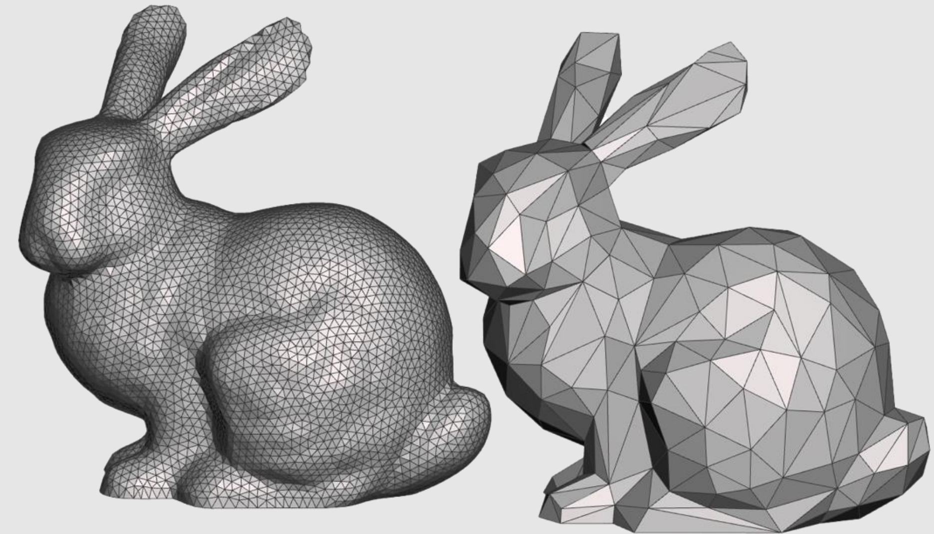
Step 2:
Flip new edges until they touch two new vertices



- ~~Good Geometry~~
- ~~Geometric Subdivision~~
- Geometric Simplification
- Geometric Remeshing
- Geometric Queries

Simplification

- Simplification is the process of **downsampling** a mesh
    - Less Storage overhead
        - Smaller file sizes
    - Less Processing overhead
        - Less elements to iterate over
    - Larger mesh modifications
        - Instead of moving tens of smaller mesh elements, move one larger mesh element

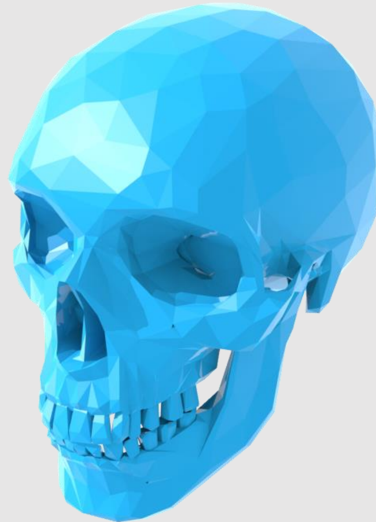


Simplification Algorithm Basics

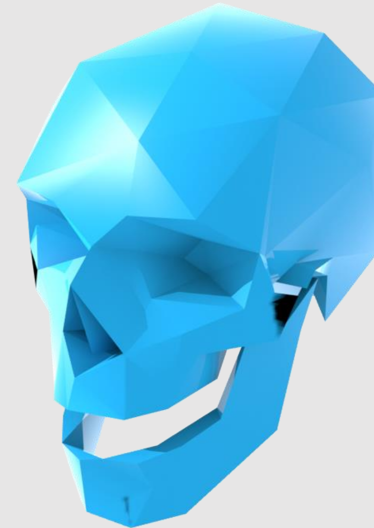
- Greedy Algorithm:
  - Assign each edge a cost
  - Collapse edge with least cost
  - Repeat until target number of elements is reached

- Particularly effective cost function: **quadric error metric****

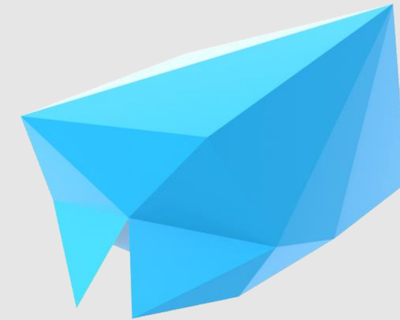
[30,000 triangles]



[3,000 triangles]



[300 triangles]



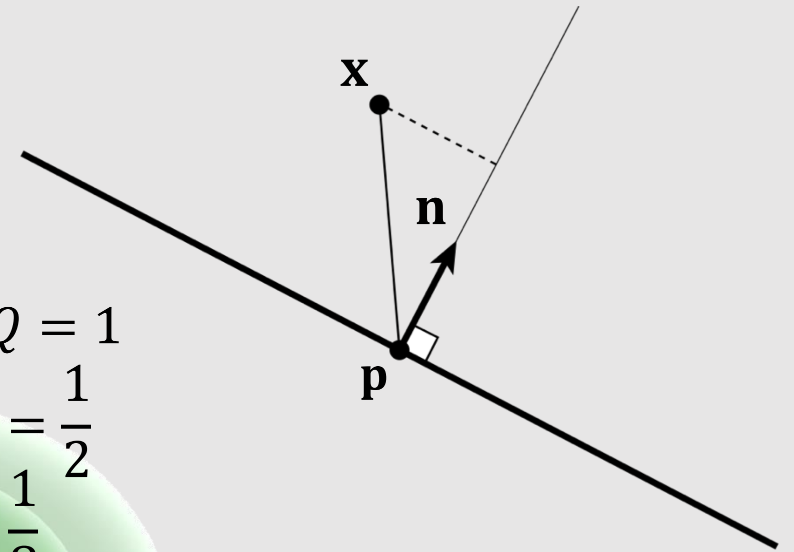
[30 triangles]

**invented at CMU (Garland & Heckbert 1997)

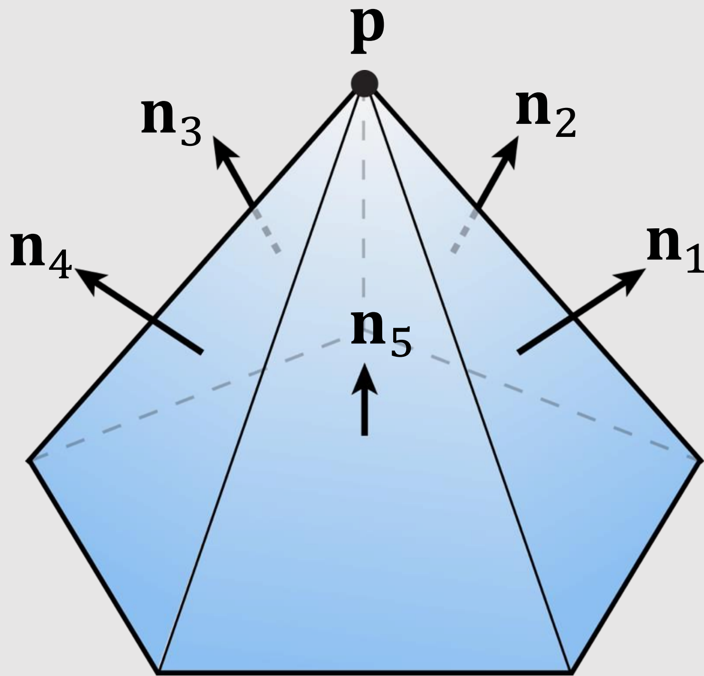
Quadratic Error Metric

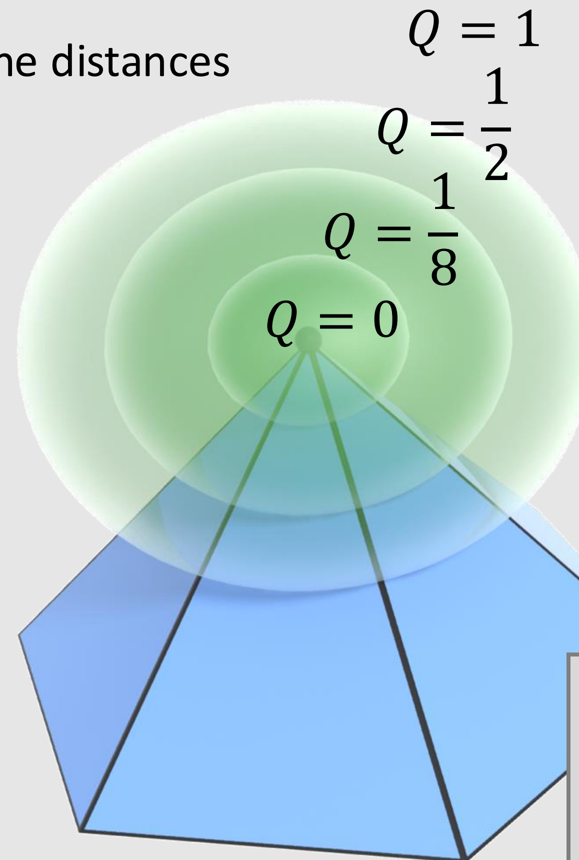
- **Goal:** approximate a point's distance from a collection of triangles
  - **Review:** what is the distance of a point $\mathbf{x}$ from a plane $\mathbf{p}$ with normal $\mathbf{n}$?

$$\mathrm{dist}(\mathbf{x}) = \langle \mathbf{n}, \mathbf{x} \rangle - \langle \mathbf{n}, \mathbf{p} \rangle = \langle \mathbf{n}, \mathbf{x} - \mathbf{p} \rangle$$



- Quadric error is the sum of squared point-to-plane distances





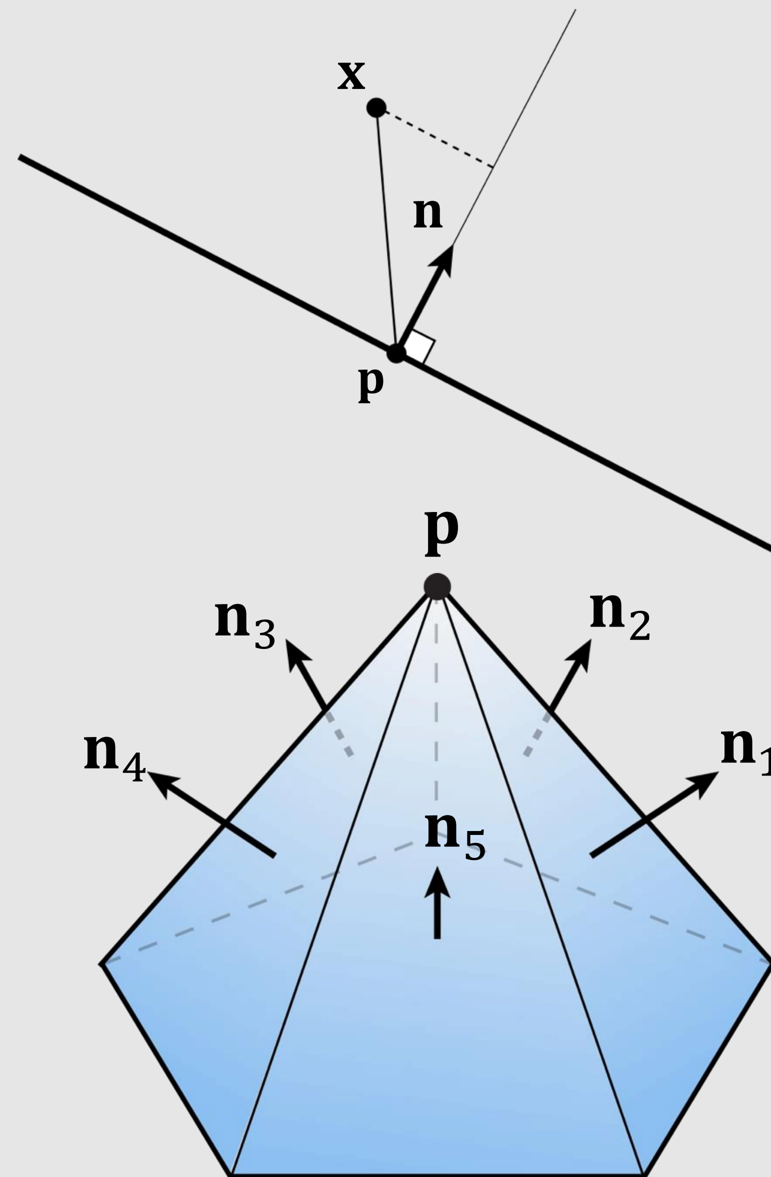
$$Q(\mathbf{x}) := \sum_{i=1}^{k} \langle \mathbf{n}_i, \mathbf{x} - \mathbf{p} \rangle^2$$

Quadratic Error As Homogeneous Coordinates

- Given:
  - Query point $\mathbf{x} = (x, y, z)$
  - Normal $\mathbf{n} = (a, b, c)$
  - Offset from origin $d = -\langle \mathbf{n}, \mathbf{p} - 0 \rangle = -\langle \mathbf{n}, \mathbf{p} \rangle$
- We can rewrite in homogeneous coordinates:
  - $\mathbf{u} = (x, y, z, 1)$
  - $\mathbf{v} = (a, b, c, d)$

- Signed distance to plane is then just $\langle \mathbf{u}, \mathbf{v} \rangle = ax + by + cz + d$
- Squared distance is $\langle \mathbf{u}, \mathbf{v} \rangle^2 = \mathbf{u}^\mathsf{T}(\mathbf{v}\mathbf{v}^\mathsf{T})\mathbf{u} =: \mathbf{u}^\mathsf{T} K \mathbf{u}$
  - Matrix $K = \mathbf{v}\mathbf{v}^T$ encodes squared distance to plane

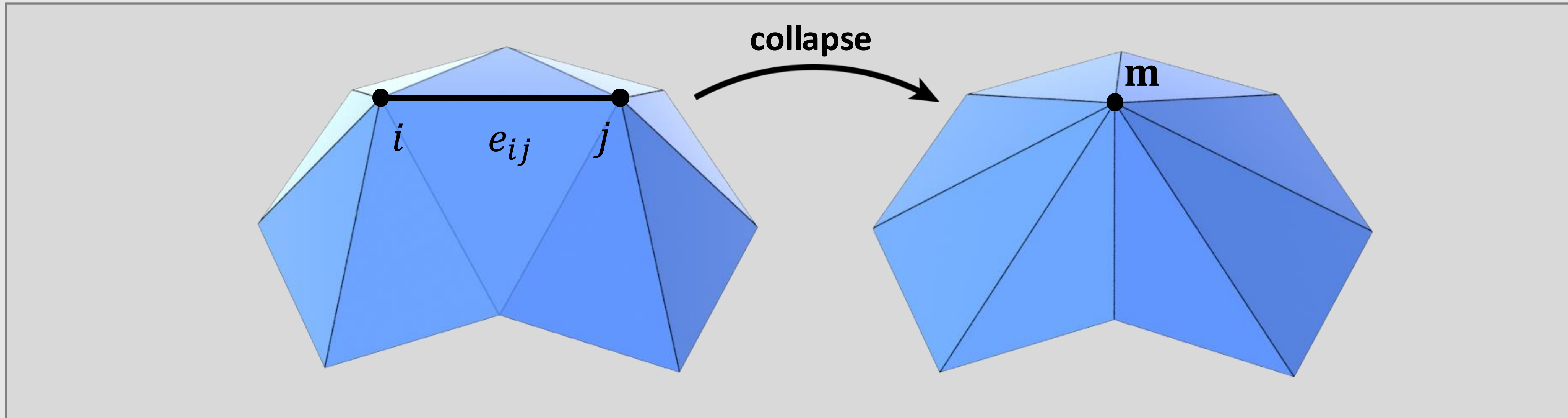
$$K = \begin{bmatrix} a^2 & ab & ac & ad \\ ab & b^2 & bc & bd \\ ac & bc & c^2 & cd \\ ad & bd & cd & d^2 \end{bmatrix}$$



Key Idea: sum of matrices K approximates distance to the mesh locally

$$\mathbf{u}^\mathsf{T} K_1 \mathbf{u} + \mathbf{u}^\mathsf{T} K_2 \mathbf{u} = \mathbf{u}^\mathsf{T}(K_1 + K_2)\mathbf{u}$$

Quadratic Error of Edge Collapse

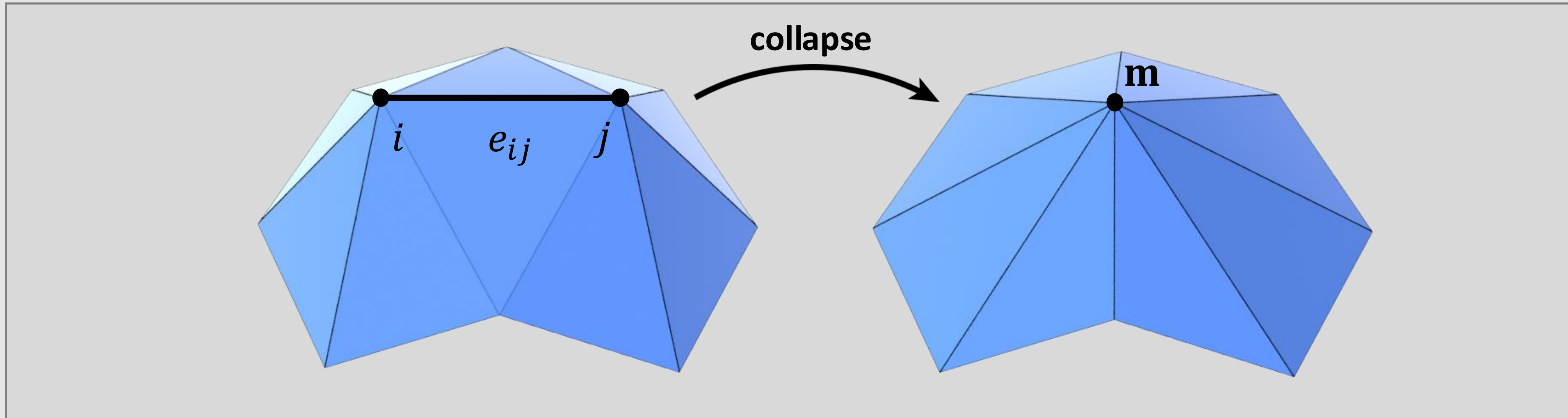


- How much does it cost to collapse an edge e_{ij}?
  - Compute midpoint $\mathbf{m}$, measure error as

$$Q(\mathbf{m}) = \mathbf{m}^\mathsf{T}(K_i + K_j)\mathbf{m}$$

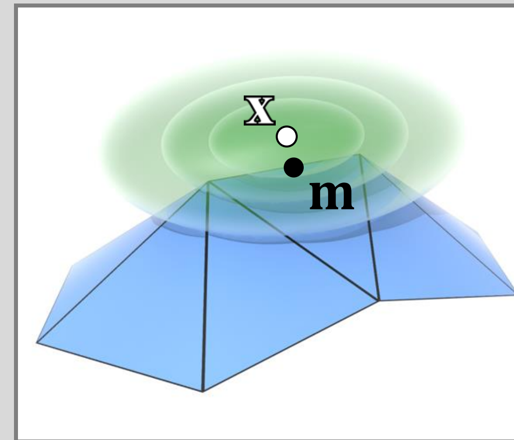
- Error becomes "score" for e_{ij}, determining priority
  - Q: where to put $\mathbf{m}$?

Quadratic Error of Edge Collapse



$$Q(\mathbf{m}) = \mathbf{m}^\mathsf{T}(K_i + K_j)\mathbf{m}$$

- Find point $\mathbf{x}$ that minimizes error
  - Take derivatives!



How to take a derivative of a function involving matrices?

Minimizing a Quadratic Function

To find the min of a function $f(x)$

$$f(x) = ax^2 + bx + c$$

take derivative $f'(x)$ and set equal to 0

$$f'(x) = 2ax + b = 0$$

$$x = -b/2a$$

same structure

can also write any quadratic function of n variables as a symmetric matrix A

consider the multivariable function

$$f(x, y) = ax^2 + bxy + cy^2 + dx + ey + g$$

we can rewrite it as:

$$\mathbf{x} = \begin{bmatrix} x \\ y \end{bmatrix} \quad A = \begin{bmatrix} a & b/2 \\ b/2 & c \end{bmatrix} \quad \mathbf{u} = \begin{bmatrix} d \\ e \end{bmatrix}$$

$$f(x, y) = \mathbf{x}^\mathsf{T} A \mathbf{x} + \mathbf{u}^\mathsf{T} \mathbf{x} + g$$

take derivative $f'(x)$ and set equal to 0

$$f'(x, y) = 2A\mathbf{x} + \mathbf{u} = 0$$

$$\mathbf{x} = -\frac{1}{2} A^{-1} \mathbf{u}$$

same structure

Positive Definite Quadratic Form

How do we know if our solution minimizes quadratic error?

$$\mathbf{x} = -\frac{1}{2}A^{-1}\mathbf{u}$$

In the 1D case, we minimize the function if

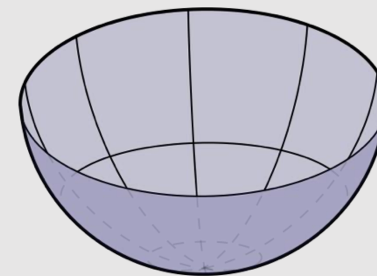
$$xax = ax^2 > 0$$
$$a > 0$$

In the N-D case, we minimize the function if

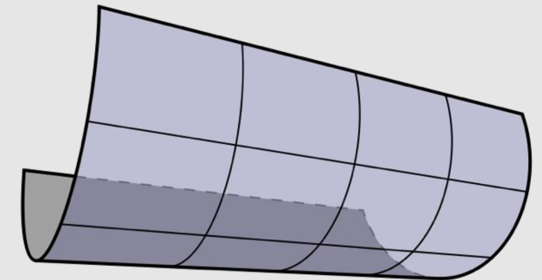
$$\boldsymbol{x}^T A \boldsymbol{x} > 0 \qquad \forall \boldsymbol{x} \neq 0$$

This is known as the function being positive-definite

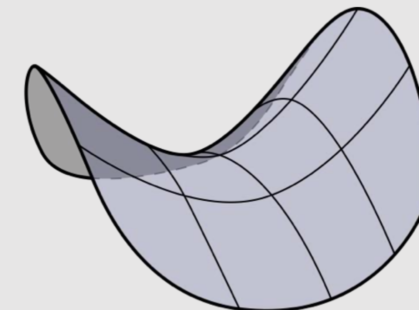
Plot of $f(\boldsymbol{x}) = \boldsymbol{x}^T A \boldsymbol{x}$ when A is …



[positive definite]



[positive semidefinite]



[indefinite]

Minimizing Quadratic Error

Find “best” point for edge collapse by minimizing quadratic form

$$\min_{\mathbf{u} \in \mathbb{R}^4} \mathbf{u}^T K \mathbf{u}$$

Already know fourth (homogeneous) coordinate for a point is 1
Break up our quadratic function into two pieces

$$\begin{bmatrix} \mathbf{x}^\mathsf{T} & 1 \end{bmatrix} \begin{bmatrix} B & \mathbf{w} \\ \mathbf{w}^\mathsf{T} & d^2 \end{bmatrix} \begin{bmatrix} \mathbf{x} \\ 1 \end{bmatrix}$$

$$= \mathbf{x}^\mathsf{T} B \mathbf{x} + 2\mathbf{w}^\mathsf{T} \mathbf{x} + d^2$$

Can minimize as before

$$2B\mathbf{x} + 2\mathbf{w} = 0$$
$$\mathbf{x} = -B^{-1}\mathbf{w}$$

Quadratic Error Simplification Algorithm

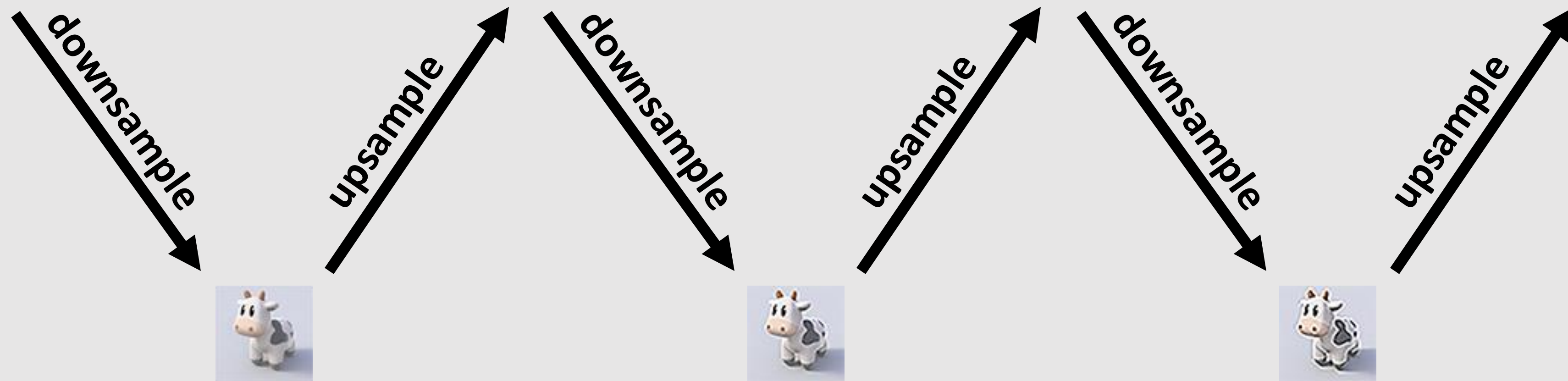
```
// compute K for each face
for(v : vertices) {
        for(f : faces) {
                Vec4 ve(N, d);
                f->K = outer(ve, ve);
        }
}

// compute K for each vertex
for(v : vertices)
        v->K.setZero()
        for(f : v->faces())
                v->K += f->K;

// compute K for each edge
// place into priority queue
PriorityQueue pq;
for(e : edge) {
        e->K.setZero()
        for(v : e->vertices())
                e->K += v->K;
        pq.push(e->K, e);
}
```

```
// iterate until mesh is a target size
while(faces.length() < target_size) {

        // collapse edge with smallest cost
        e = pq.pop();
        K = e->K;
        v = collapse(e);

        // position new vertex to optimal pos
        v->pos = -B.inv() * w

        // update K for vertex
        // update K for edges touching vertex
        v->K = K;
        for(e2 : v->edges()) {
                e2->K.setZero()
                for(v2 : e2->vertices())
                        e2->K += v2->K;
        }
}
```

Note: we are not recomputing face quadrics, because we want to stay close to the **original** mesh.

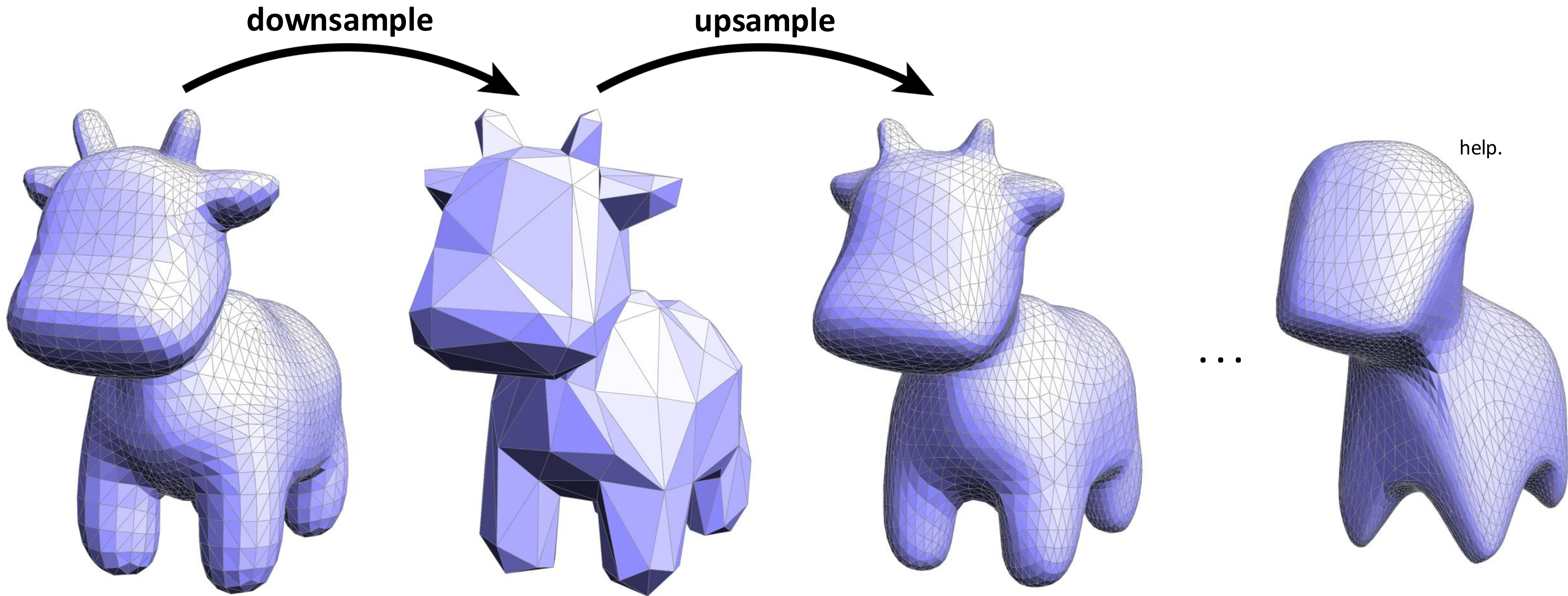
Is simplification the inverse operation of subdivision?

Dangers of Resampling



Repeatedly resampling an image degrades signal quality!

Dangers of Resampling



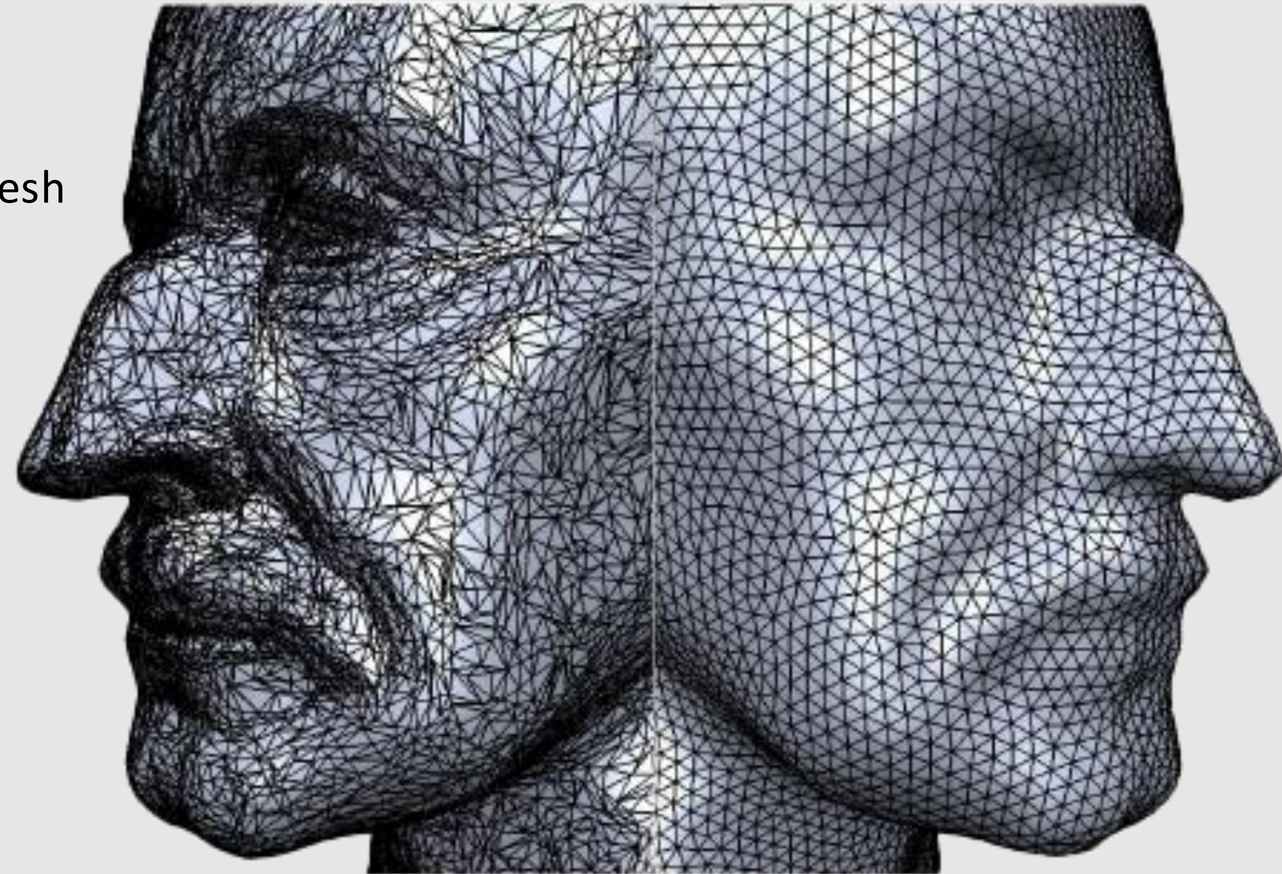
Repeatedly downsampling and upsampling a mesh also degrades signal quality!

- ~~Good Geometry~~
- ~~Geometric Subdivision~~
- ~~Geometric Simplification~~
- Geometric Remeshing
- Geometric Queries

Isotropic Remeshing

- **Isotropic:** same value when measured in any direction
- **Remeshing:** a change (often inc. connectivity) in the mesh
    - **Goal:** change the mesh to make triangles more uniform shape and size

- Helps achieve good mesh properties:
    - Good approximation of original shape
    - Vertex degrees close to 6
    - Angles close to 60deg
    - Delaunay triangles

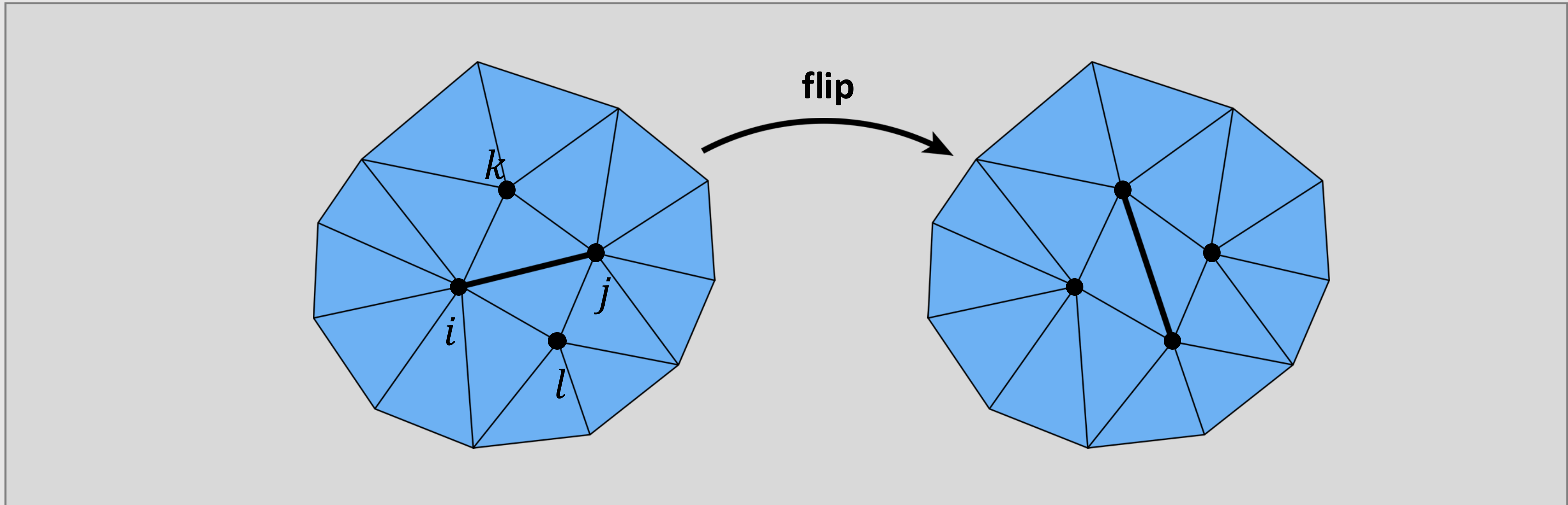


Improving Degree

Vertices with degree 6 makes triangles more regular

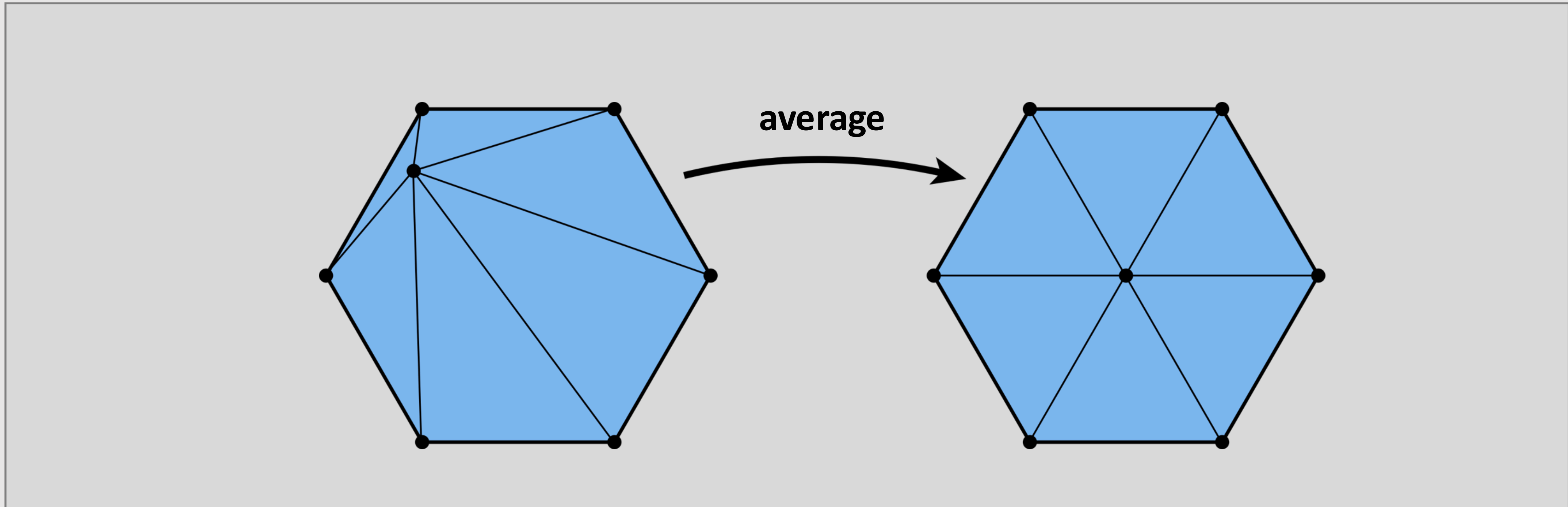
Deviation function: $|d_i - 6| + |d_j - 6| + |d_k - 6| + |d_l - 6|$

If flipping an edge reduces deviation function, flip edge



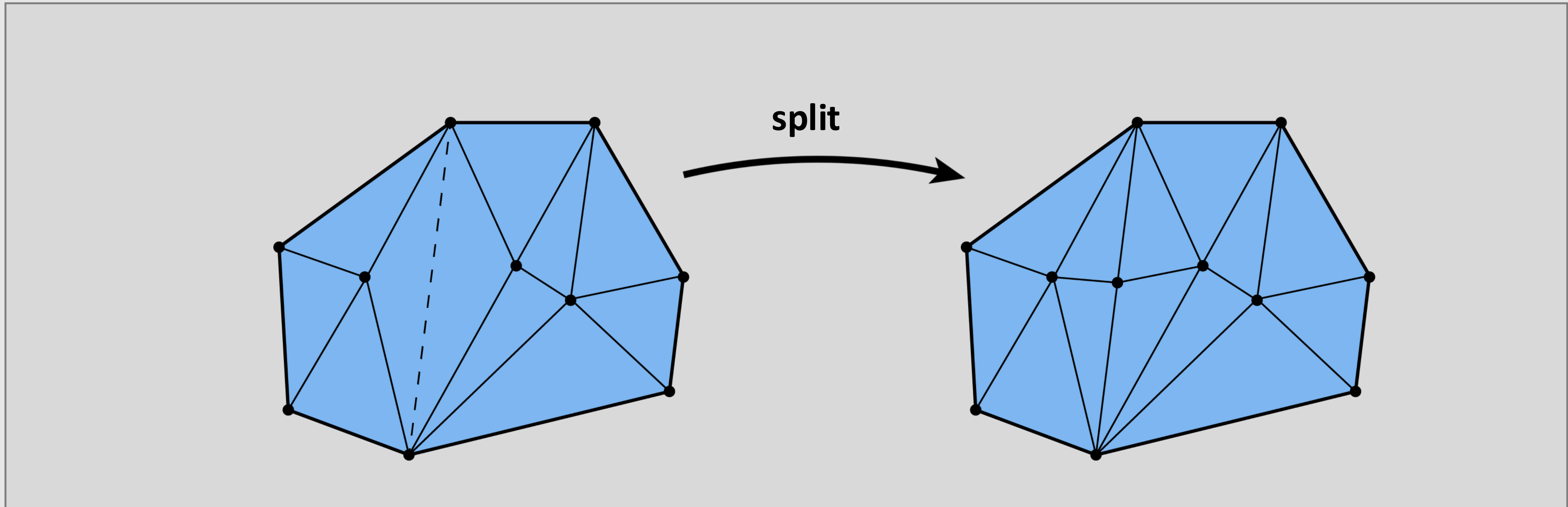
Improving Vertex Positioning

Center vertices to make triangles more even in size



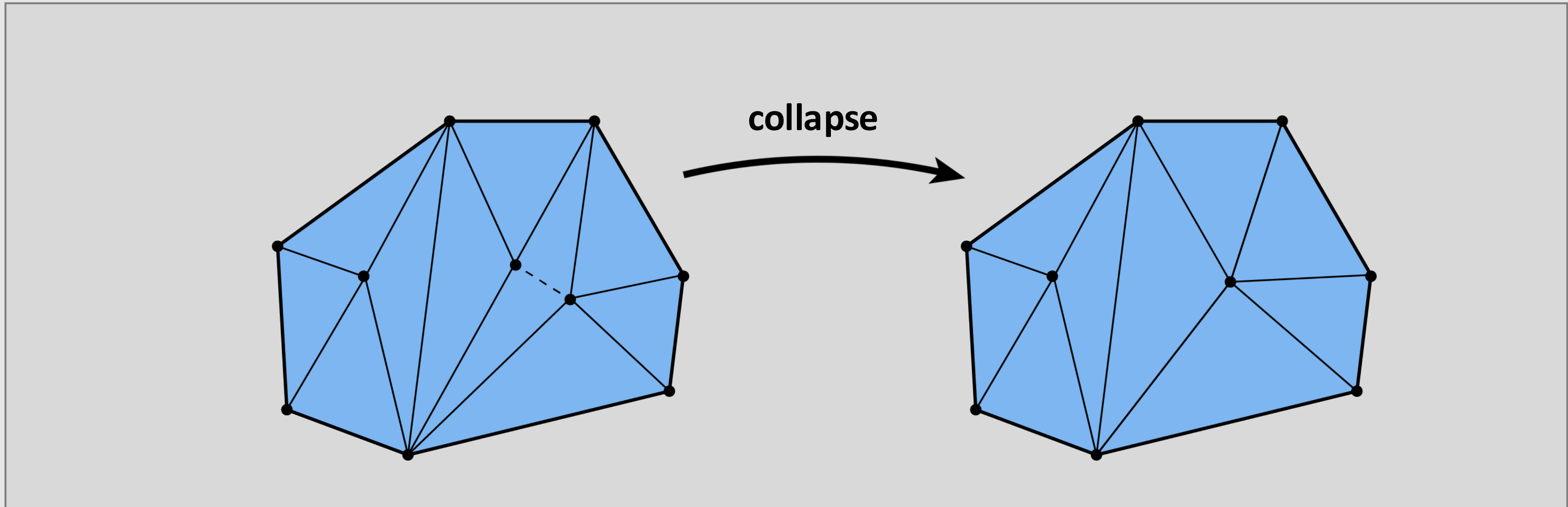
Improving Edge Length

If an edge is longer than (4/3 * mean) length, split it

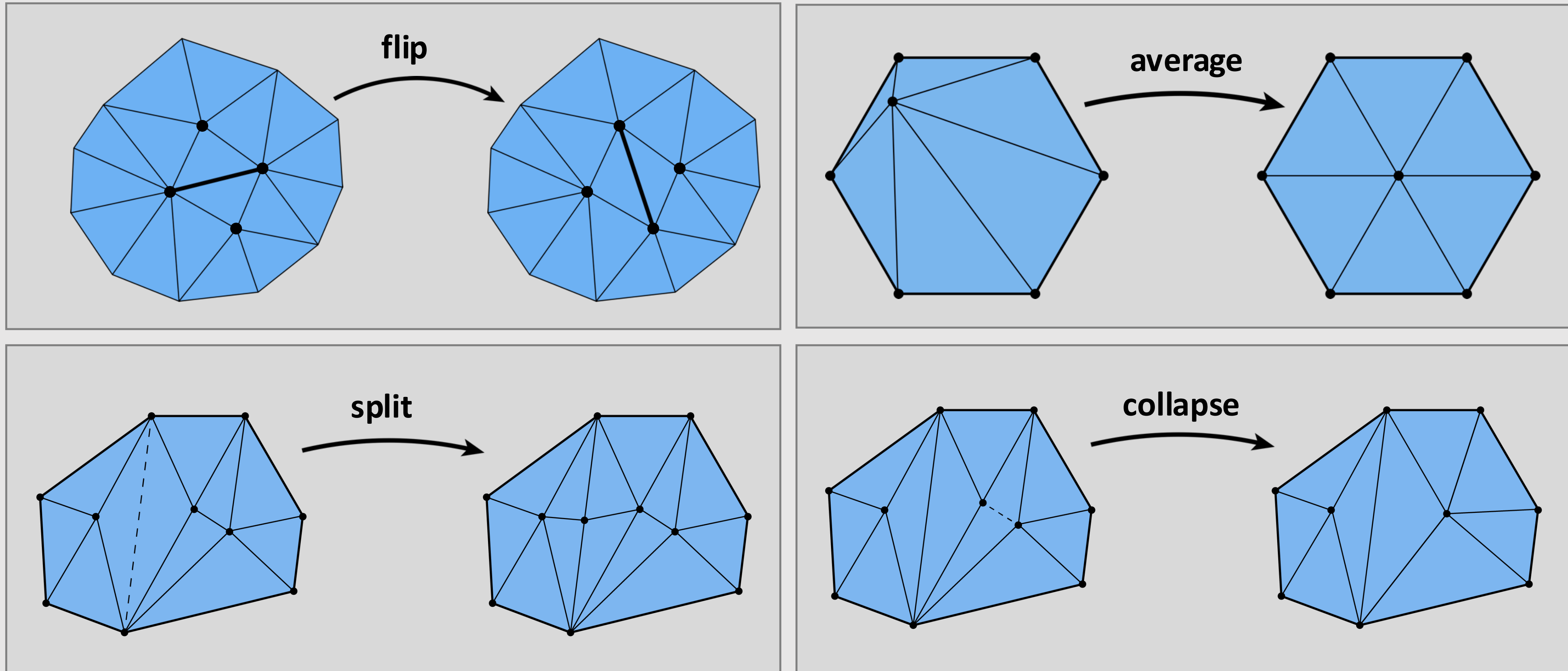


Improving Edge Length

If an edge is shorter than (4/5 * mean) length, collapse it



Isotropic Remeshing

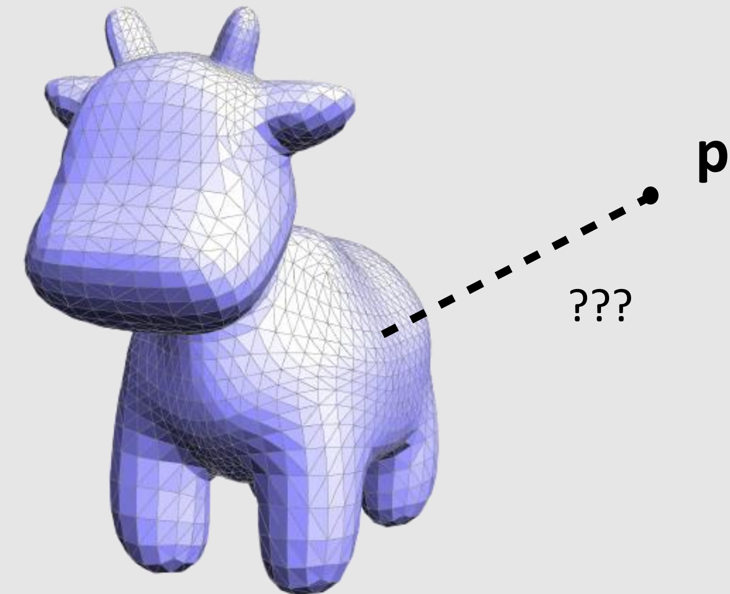


Note: all these operations may significantly change the mesh geometry locally, esp. on coarse meshes.
Extra consideration can be taken to decide new vertex position and whether this operation should be performed or not. (Extra consideration not required in the homework, except for averaging.)

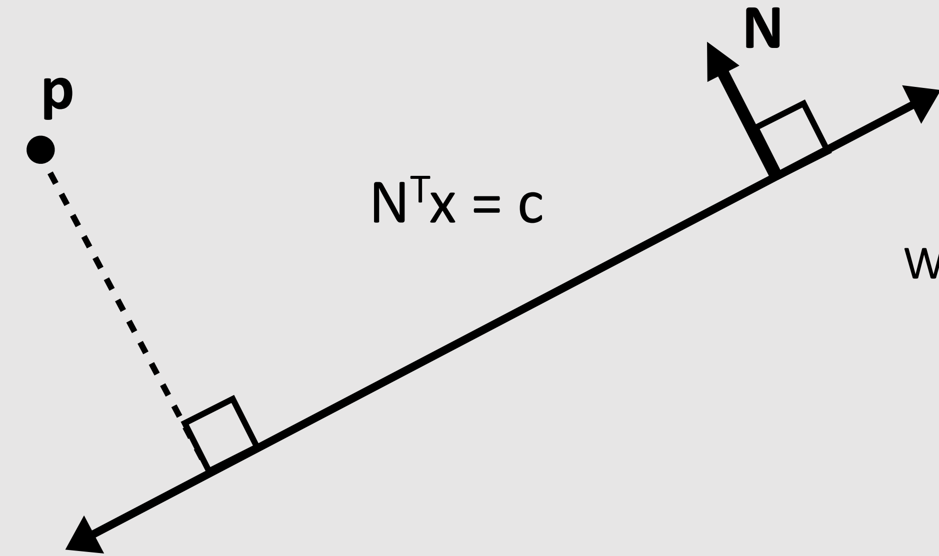
- ~~Good Geometry~~
- ~~Geometric Subdivision~~
- ~~Geometric Simplification~~
- ~~Geometric Remeshing~~
- Geometric Queries

Closest Point Queries

- **Problem:** given a point, how do we find the closest point on a given surface?
- Several use cases:
  - Ray/mesh intersection in path tracing
  - Kinematics/animation
  - GUI/user selection
    - When I click on a mesh, what point am I actually clicking on?



Closest Point on a Line



To find the closest point to **p** along $N^T x = c$
We can have **p** travel along **N** for some time t
to arrive at the closest point:

$$N^T(p + tN) = c$$

Multiplying the terms out

$$N^T p + tN^T N = c$$

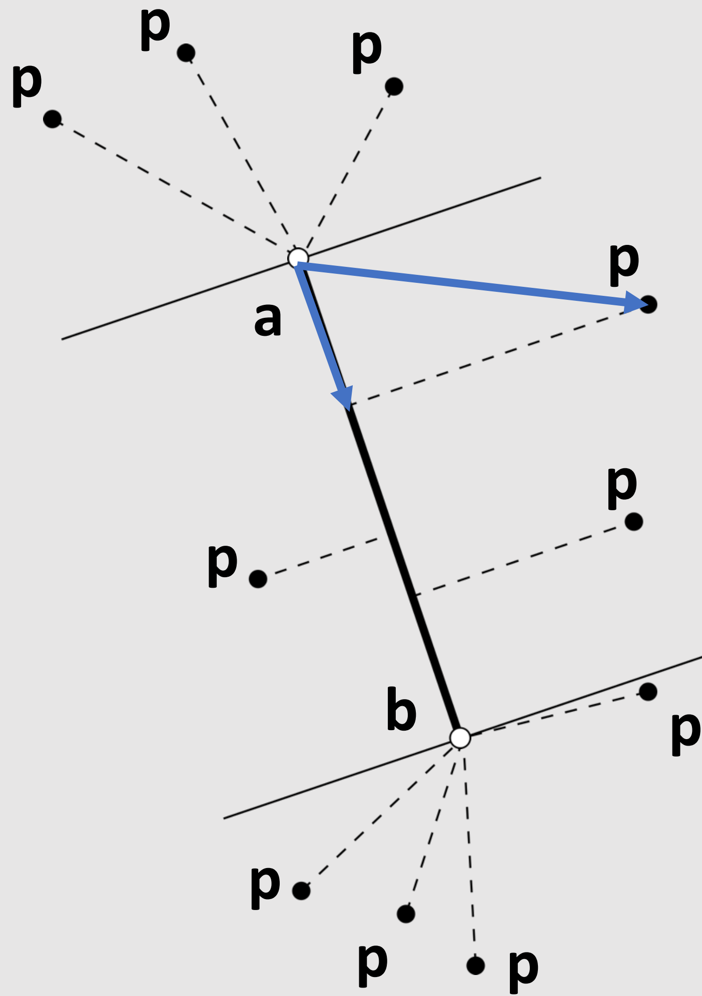
The unit norm multiplied by itself is 1
Solve for t

$$t = c - N^T p$$

Propagate **p** along **N** for time t

$$p + tN$$
$$p + (c - N^T p)N$$

Closest Point on a Line Segment



Compute the projection of **ap** onto the line of **ab**

$$\langle \mathbf{p} - \mathbf{a}, \mathbf{b} - \mathbf{a} \rangle$$

Normalize to get a time (traveling from **a** to the closest point to the line)

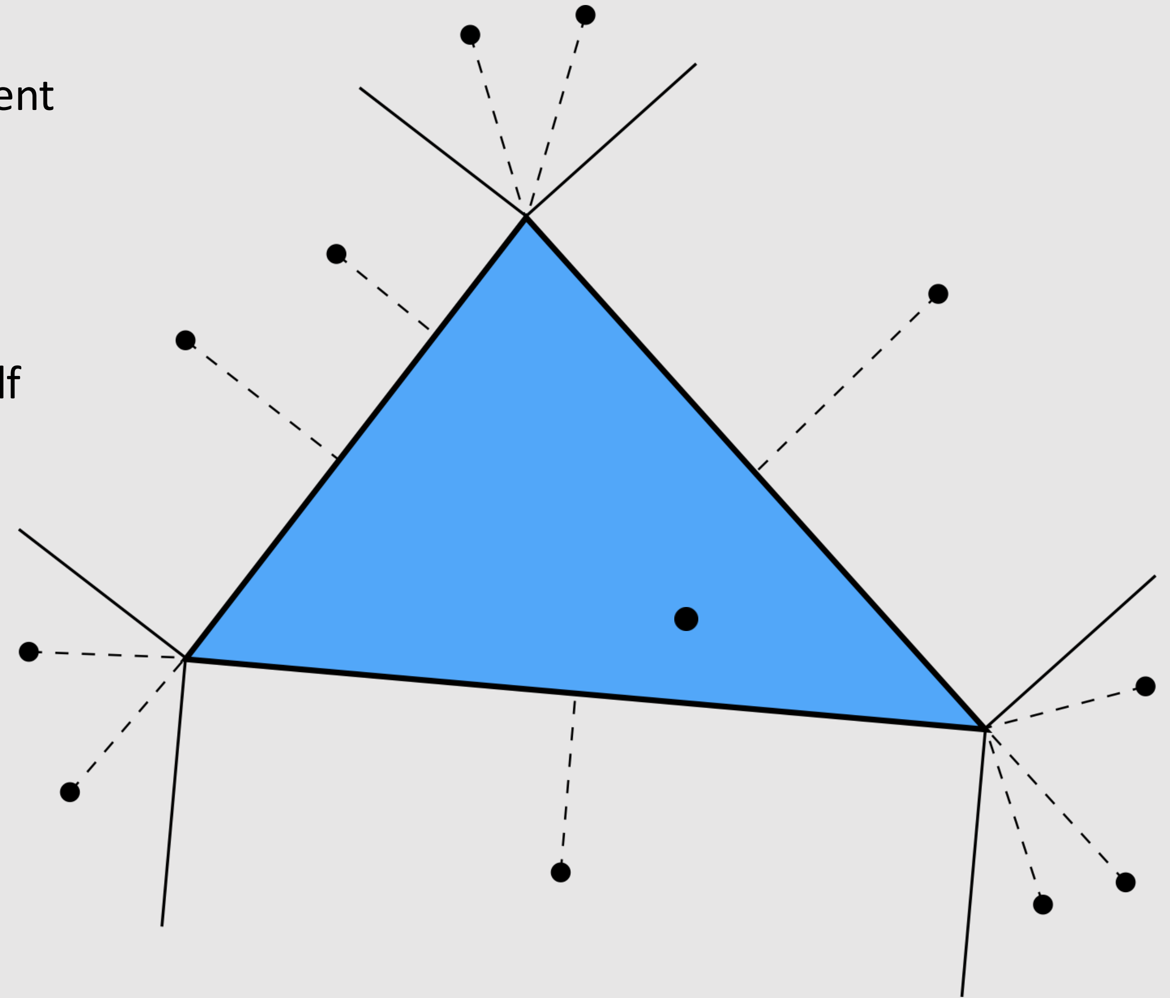
$$t = \frac{\langle \mathbf{p} - \mathbf{a}, \mathbf{b} - \mathbf{a} \rangle}{\langle \mathbf{b} - \mathbf{a}, \mathbf{b} - \mathbf{a} \rangle}$$

Clip time to range [0,1] and get the closet point to the segment

$$\boldsymbol{a} + (\mathbf{b} - \mathbf{a})t$$

Closest Point on a 2D Triangle

- Easy! Just compute closest point to each line segment
  - For each point, compute distance
  - Point with smallest distance wins
- What if the point is inside the triangle?
  - Even easier! The closest point is the point itself
  - Recall point-in-triangle tests



Closest Point on a 3D Triangle

- **Method #1:** Projection**
  - Construct a plane that passes through the triangle
    - Can be done using cross product of edges
  - Project the point to the closest point on the plane
    - Same expression as with a line: $p + (c - N^T p)N$
    - Check if point is in triangle using half-plane test
  - Else, compute distance from each line segment in 3D
    - Same expression as with a 2D line segment

- **Method #2:** Rotation**
  - Translate point + triangle so that triangle vertex v1 is at the origin
  - Rotate point + triangle so that triangle vertex v2 sits on the z-axis
  - Rotate point + triangle so that triangle vertex v3 sits on the yz-axis
  - Disregard x-coordinate of point
    - Problem reduces to closest point on 2D triangle

**https://citeseerx.ist.psu.edu/viewdoc/download?doi=10.1.1.104.4264&rep=rep1&type=pdf

Closest Point on a 3D Triangle Mesh

- Conceptually easy!
  - Loop over every triangle
  - Compute closest point to current triangle
  - Keep track of globally closest point

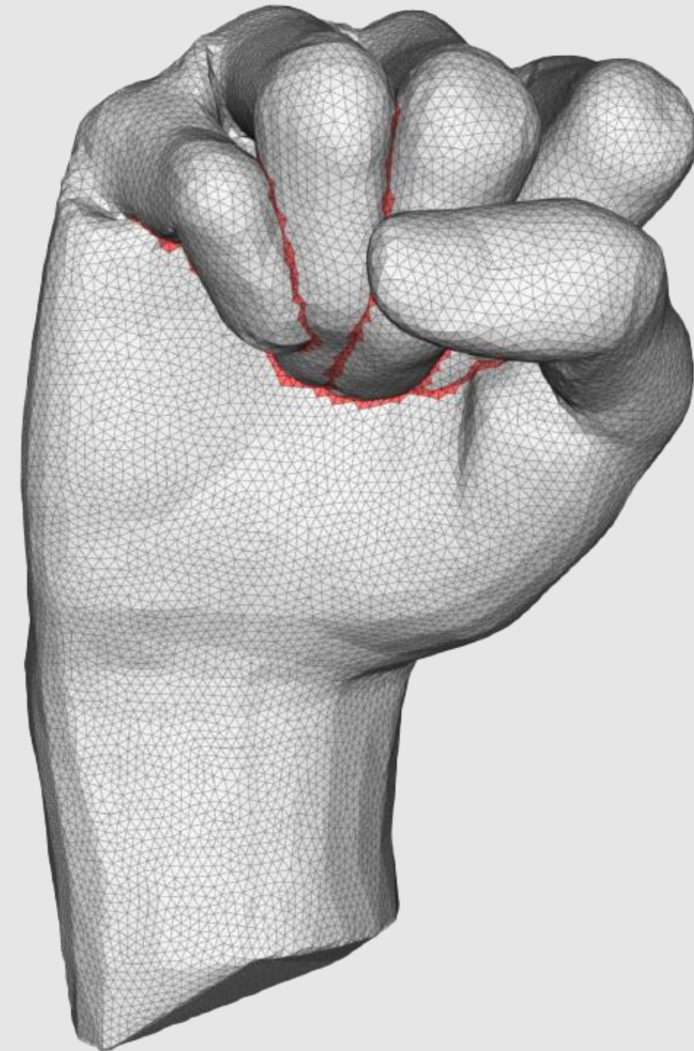
- Not practical in real world
  - Meshes can have millions of triangles
  - Programs make thousands of geometric queries a second

- Will look at better solutions next time

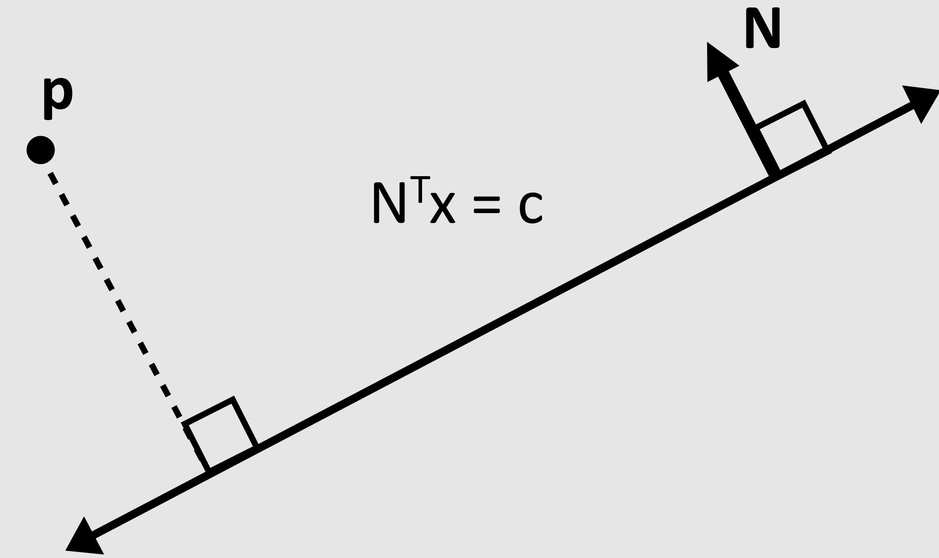
Mesh-Mesh Intersections

- Sometimes when editing geometry, a mesh will intersect with itself
- Likewise, sometimes when animating geometry, meshes will collide

- How do we check for/prevent intersections?



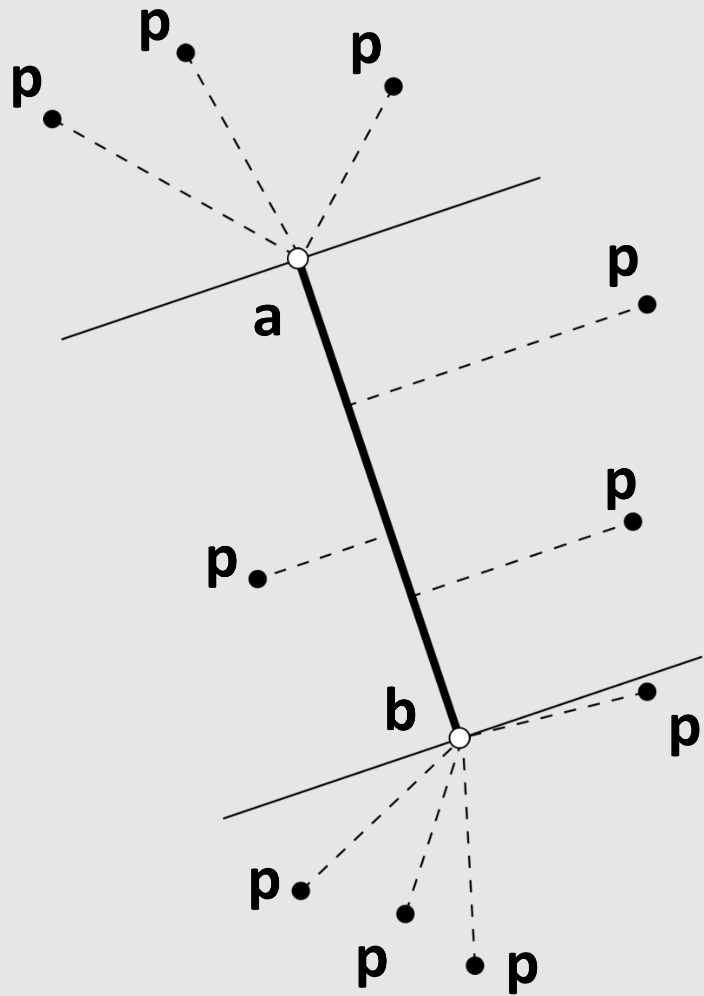
Point-Line Intersection



Just plug point in

$$N^T p = c?$$

Point-Line Segment Intersection



Check if adding distances equals net distance**

$$dist(a,p) + dist(p,b) = dist(a,b)$$

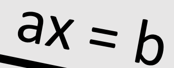
**Potential numeric stability issues

Line-Line Intersection

Two equations, two unknowns

Solve a linear system

$$\begin{bmatrix} a_1 & a_2 \\ c_1 & c_2 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} b \\ d \end{bmatrix}$$



cx = d

Point-Triangle Intersection

You know this :)